

Approximate planning and learning for partially observed systems

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Joint work with Jayakumar Subramanian
Thanks to Amit Sinha and Raihan Seraj for simulation results

Mila RL Reading Group
14 February 2020

Many successes of RL in recent years

- ▷ Algorithms based on comprehensive theory

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Alpha Go

Approx. POMDPs-(Mahajan)

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Arcade games

Approx. POMDPs-(Mahajan)

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Robotics

Approx. POMDPs-(Mnih et al.)



Robotics

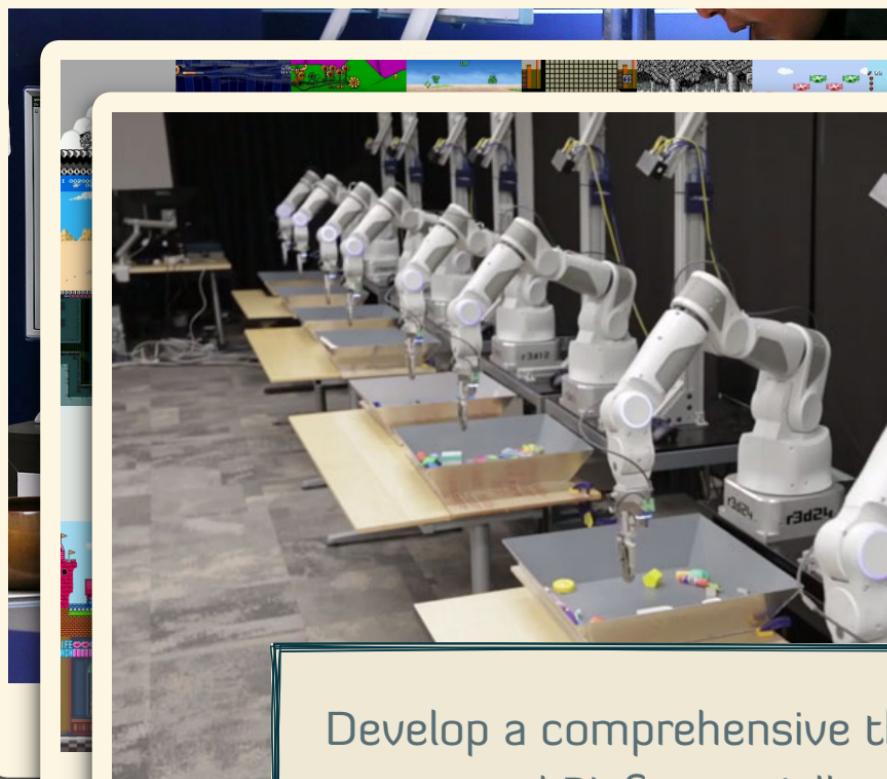
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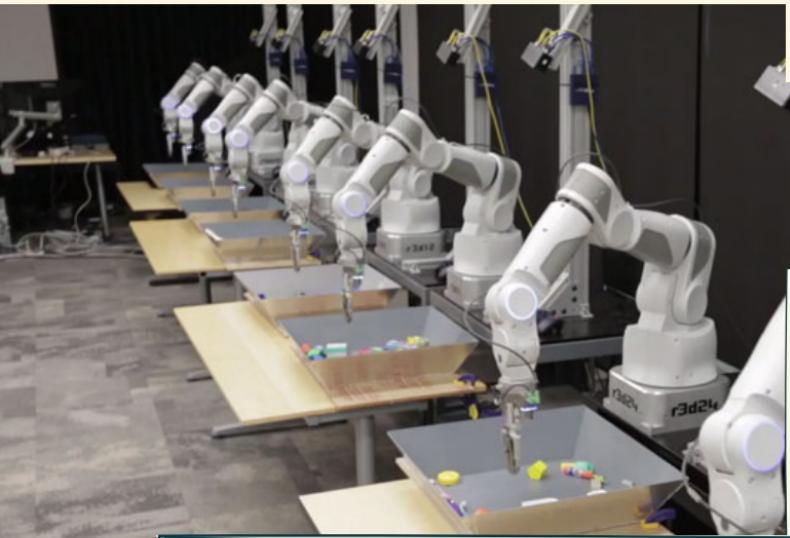
Applications with partially observed state

- ▷ Healthcare
- ▷ Autonomous driving
- ▷ Finance (portfolio management)
- ▷ Retail and marketing



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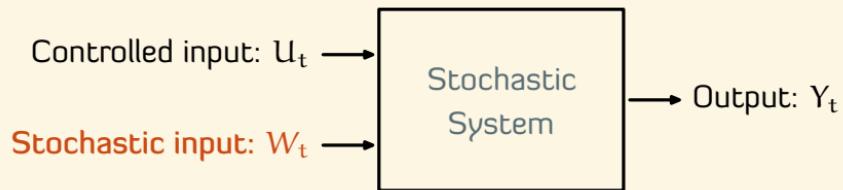
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Develop a comprehensive theory of approximate DP
and RL for partially observed systems

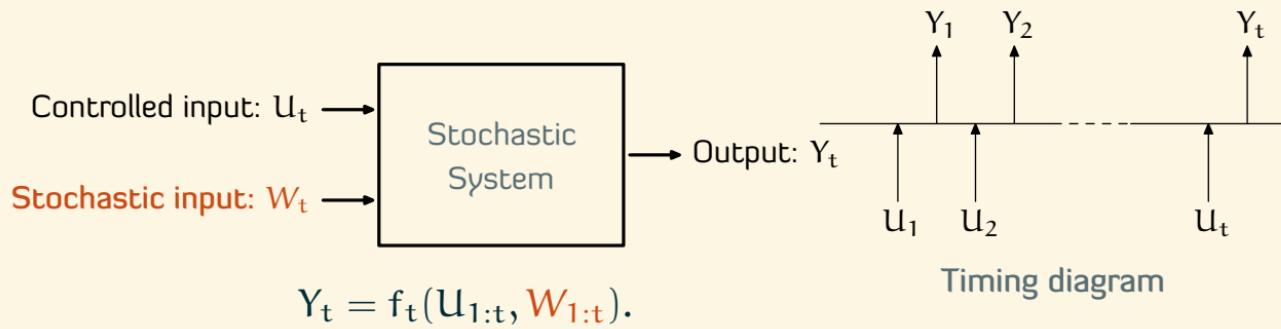
Belief state for partially observed systems

Belief state in partially observed stochastic dynamical systems

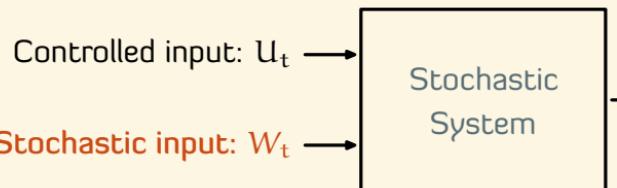


$$Y_t = f_t(U_{1:t}, W_{1:t}).$$

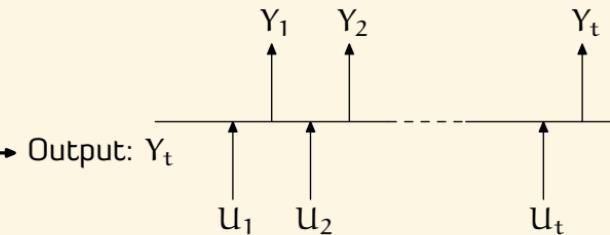
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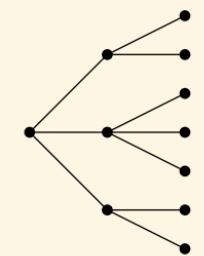
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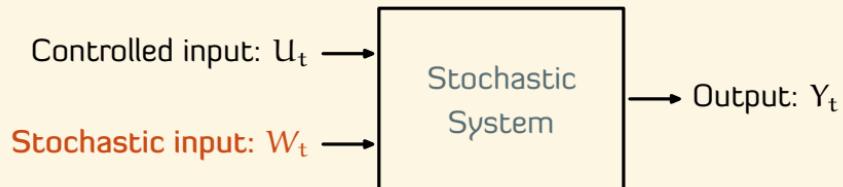
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Timing diagram



Belief state in partially observed stochastic dynamical systems

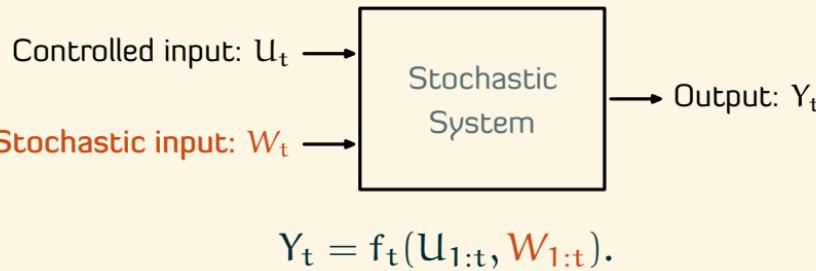


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STOCHASTIC INPUT IS NOT OBSERVED

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Belief state in partially observed stochastic dynamical systems



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TRADITIONAL SOLUTION: BELIEF STATES

Step 1 Identify a state $\{S_t\}_{t \geq 0}$ for predicting output assuming that the stochastic inputs are observed.

Step 2 Define a BELIEF STATE $B_t \in \Delta(\mathcal{S})$:

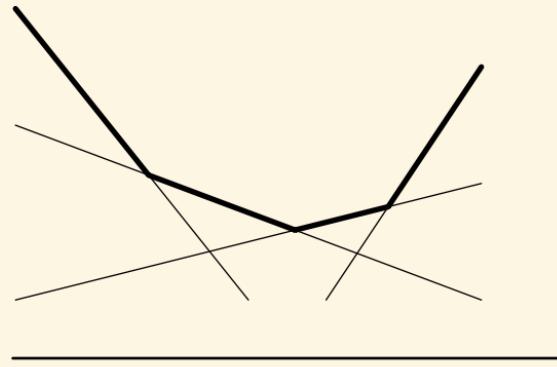
$$B_t(s) = \mathbb{P}(S_t = s | H_t = h_t), \quad s \in \mathcal{S}.$$

► Astrom, "Optimal control of Markov decision processes with incomplete state information," 1965. ► Striebel, "Sufficient statistics in the optimal control of stochastic systems," 1965.

► Stratonovich, "Conditional Markov processes," 1960. ► Baum and Petrie, "Statistical inference for probabilistic functions of finite state Markov chains," 1966.

Partially observed Markov decision processes (POMDPs): Pros and Cons of belief state representation

Value function is piecewise linear and convex.



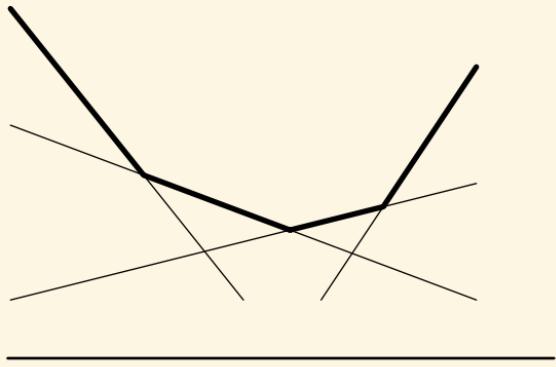
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Approx. POMDPs-(Mahajan)

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When the state space model is not known analytically (as is the case for black-box models and simulators as well as some real world application such as healthcare), belief states are difficult to construct and difficult to approximate from data.

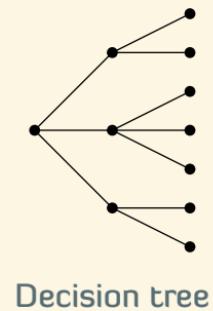
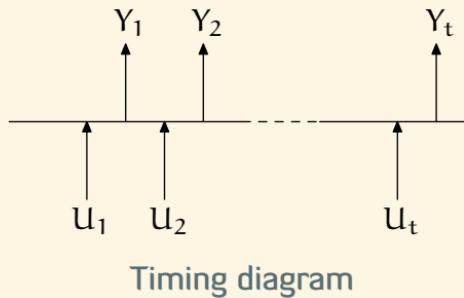
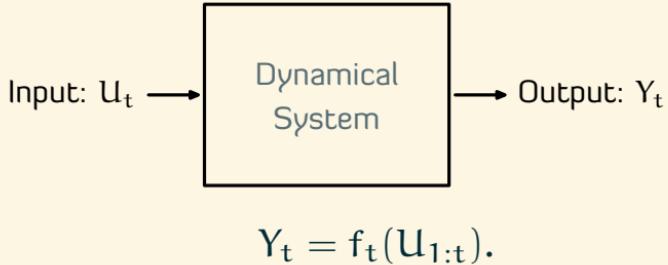
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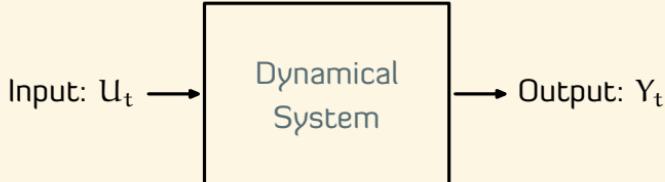
Is there another ways to model
partially observed systems which is
more amenable to approximations?

Let's go back to first principles.

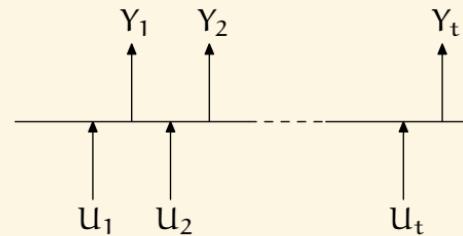
Notion of state in deterministic dynamical systems



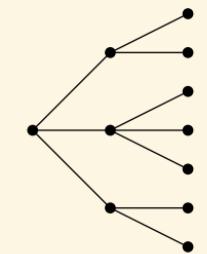
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Timing diagram



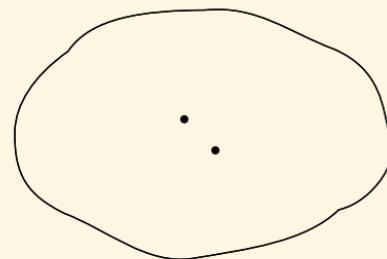
Decision tree

EQUIVALENCE RELATIONSHIP

Let $H_t = U_{1:t-1}$ denote the history of inputs until time t .

$H_t^{(1)} \sim H_t^{(2)}$ if for all future inputs $U_{t:T}$, the future outputs $Y_{t:T}^{(1)}$ and $Y_{t:T}^{(2)}$ are the same:

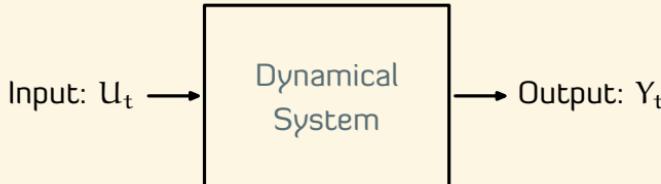
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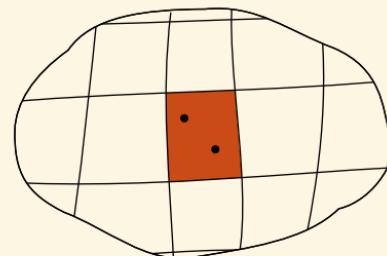
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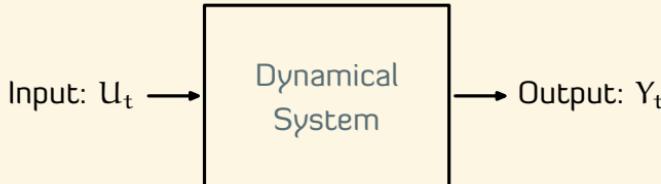
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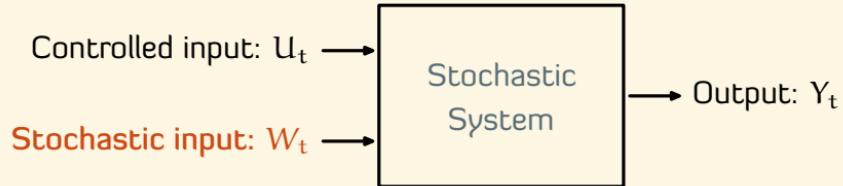
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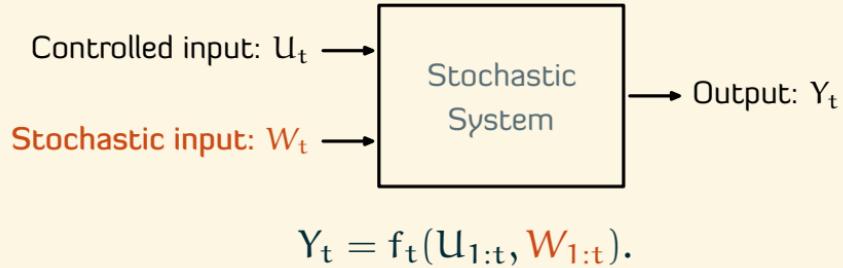
(Ignore: measurability and minimality)

Notion of state in stochastic dynamical systems



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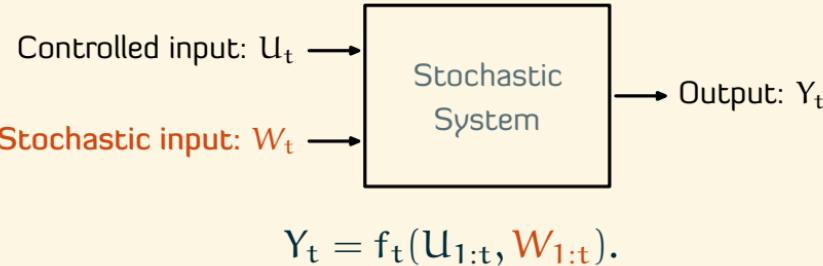


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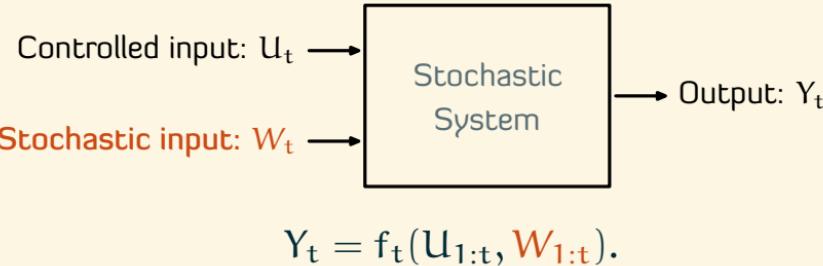
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- ▷ Kalman, “Mathematical description of linear dynamical systems”, 1963.
- ▷ Balakrishnan, “Foundations of state-space theory of cts systems”, 1967.
- ▷ Willems, “The generation of Lyapunov functions for I/O stable systems”, 1971.

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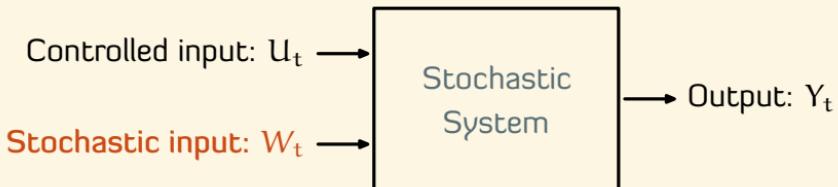
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We recover the two basic models
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What happens when the
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Notion of state in partially observed stochastic dynamical systems

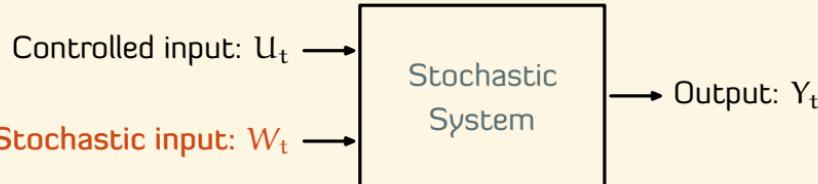


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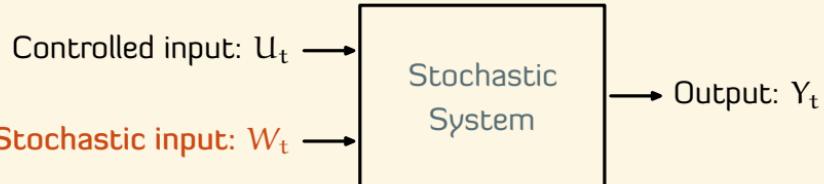
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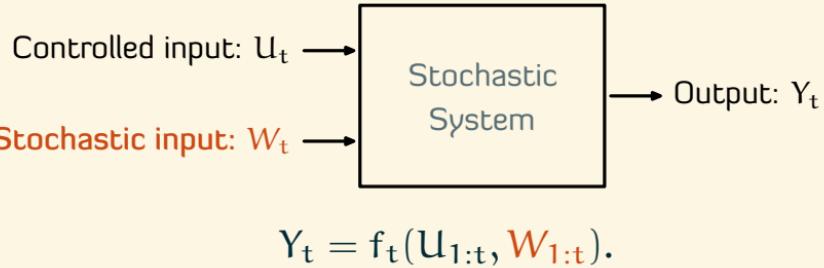
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Too restrictive . . .

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PROPERTIES OF INFORMATION STATE

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▷ SUFFICIENT TO PREDICT OUTPUT:

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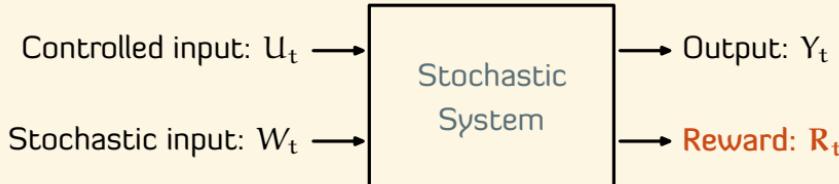
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KEY QUESTIONS

- ▷ Can this be used for dynamic programming?
- ▷ What is the right notion of approximations in this framework?

An information state for dynamic programming

Predicting output vs optimizing expected rewards over time

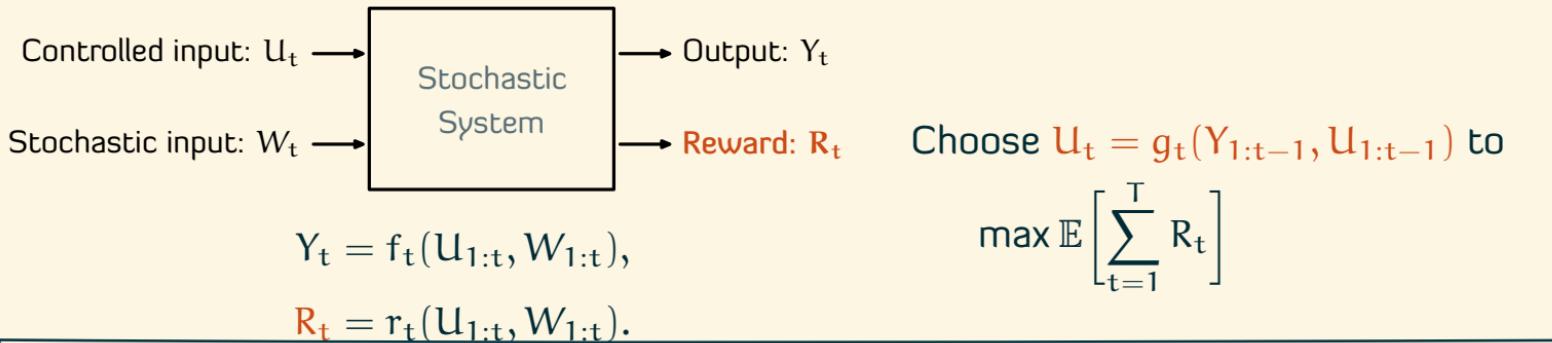


$$Y_t = f_t(U_{1:t}, W_{1:t}),$$
$$R_t = r_t(U_{1:t}, W_{1:t}).$$

Choose $U_t = g_t(Y_{1:t-1}, U_{1:t-1})$ to

$$\max \mathbb{E} \left[\sum_{t=1}^T R_t \right]$$

Predicting output vs optimizing expected rewards over time



PROPERTIES OF INFORMATION STATE (SUFFICIENT FOR DYNAMIC PROGRAMMING)

The info state Z_t at time t is a “compression” of past inputs that satisfies the following:

▷ SUFFICIENT TO PREDICT ITSELF:

$$\mathbb{P}(Z_{t+1} | H_t, U_t) = \mathbb{P}(Z_{t+1} | Z_t, U_t).$$

▷ SUFFICIENT TO ESTIMATE EXPECTED REWARD:

$$\mathbb{E}[R_t | H_t, U_t] = \mathbb{E}[R_t | Z_t, U_t].$$

Dynamic programming using information state

PRELIMINARY THEOREM

If $\{Z_t\}_{t \geq 1}$ is any information state process. Then:

- ▷ There is no loss of optimality in restricting attention to policies of the form
 $u_t = \tilde{g}_t(Z_t)$.

Dynamic programming using information state

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If $\{Z_t\}_{t \geq 1}$ is any information state process. Then:

- ▷ There is no loss of optimality in restricting attention to policies of the form

$$U_t = \tilde{g}_t(Z_t).$$

- ▷ Let $\{V_t\}_{t=1}^{T+1}$ denote the solution to the following dynamic program: $V_{T+1}(z_{T+1}) = 0$ and for $t \in \{T, \dots, 1\}$,

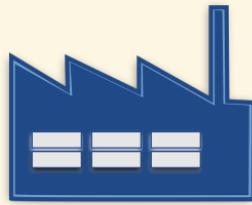
$$Q_t(z_t, u_t) = \mathbb{E}[R_t + V_{t+1}(Z_{t+1}) \mid Z_t = z_t, U_t = u_t],$$

$$V_t(z_t) = \max_{u_t \in \mathcal{U}} Q_t(z_t, u_t).$$

A policy $\{\tilde{g}_t\}_{t=1}^T$, $\tilde{g}_t: \mathcal{Z}_t \rightarrow \mathcal{U}$, is optimal if it satisfies

$$\tilde{g}_t(z_t) \in \arg \max_{u_t \in \mathcal{U}} Q_t(z_t, u_t).$$

An example: Machine repair

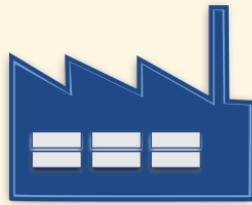


$$P(\text{RUN}) = P_0, \quad P(\text{INSPECT}) = I$$

$$P(\text{REPLACE}) = \begin{bmatrix} 1 & 0 & 0 \\ 1 & 0 & 0 \\ 1 & 0 & 0 \end{bmatrix}$$

- ▷ State $\in \{1, 2, \dots, n\}$
- ▷ Action $\in \{\text{RUN}, \text{INSPECT} + \text{REPAIR}\}$.
- ▷ $\text{cost(state,action)} = \text{running cost(state)} + \text{inspection cost} + \text{repair cost}$

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Belief state: $P(S_t | Y_{1:t-1}) \in \mathbb{R}^n$

Alternative information state $(S_\tau, t - \tau) \in \{1, \dots, n\} \times \mathbb{N}$

What about approximations?

Preliminary: A family of pseudometrics on probability distribution

INTEGRAL PROBABILITY METRIC (IPM)

Let \mathcal{P} denote the set of probability measures on a measurable space $(\mathcal{X}, \mathcal{G})$.

Given a class \mathfrak{F} of real-valued bounded measurable functions on $(\mathcal{X}, \mathcal{G})$, the integral probability metric (IPM) between two probability distributions $\mu, \nu \in \mathcal{P}$ is given by:

$$d_{\mathfrak{F}}(\mu, \nu) = \sup_{f \in \mathfrak{F}} \left| \int_{\mathcal{X}} f d\mu - \int_{\mathcal{X}} f d\nu \right|.$$

► Müller, "Integral probability metrics and their generating classes of functions," 1997.

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EXAMPLES

- ▷ If $\mathfrak{F} = \{f : \|f\|_{\infty} \leq 1\}$,
 $d_{\mathfrak{F}}$ = Total variation distance.
- ▷ If $\mathfrak{F} = \{f : |f|_L \leq 1\}$,
 $d_{\mathfrak{F}}$ = Wasserstein distance.
- ▷ If $\mathfrak{F} = \{f : \|df/dx\|_{1-1/p} \leq 1\}$,
 $d_{\mathfrak{F}}$ = Cramér p distance
- ▷ ...

We say a function f has a \mathfrak{F} -constant K if $f/K \in \mathfrak{F}$.

▷ Müller, "Integral probability metrics and their generating classes of functions," 1997.

Approximate information state

(ε, δ) -APPROXIMATE INFORMATION STATE (AIS)

Given a function class \mathfrak{F} , a compression $\{Z_t\}_{t \geq 1}$ of history (i.e., $Z_t = \varphi_t(H_t)$) is called an $\{(\varepsilon_t, \delta_t)\}_{t \geq 1}$ AIS if there exist:

- ▷ a function $\tilde{R}_t(Z_t, U_t)$, and ▷ a stochastic kernel $v_t(Z_{t+1}|Z_t, U_t)$

such that

- ▷ $|\mathbb{E}[R_t|H_t = h_t, U_t = u_t] - \tilde{R}_t(\varphi_t(h_t), u_t)| \leq \varepsilon_t$
- ▷ For any Borel set A of \mathcal{Z}_{t+1} , define

$$\mu_t(A) = \mathbb{P}(Z_{t+1} \in A | H_t = h_t, U_t = u_t)$$

Then,

$$d_{\mathfrak{F}}(\mu_t, v_t(\cdot | \varphi_t(h_t), u_t)) \leq \delta_t.$$

Approximate dynamic programming using AIS

MAIN THEOREM

Given a function class \mathcal{F} , let $\{Z_t\}_{t \geq 1}$, where

$Z_t = \varphi_t(H_t)$, be an $\{(\varepsilon_t, \delta_t)\}_{t \geq 1}$ AIS.

Recursively define the following functions:

$$\hat{V}_{T+1}(z_{T+1}) = 0$$

and for $t \in \{T, \dots, 1\}$:

$$\begin{aligned}\hat{V}_t(z_t) = \max_{u_t \in \mathcal{U}} & \left\{ \tilde{R}_t(z_t, u_t) \right. \\ & \left. + \int \hat{V}_{t+1}(z_{t+1}) \nu_t(dz_{t+1} | z_t, u_t) \right\}.\end{aligned}$$

Let $\pi = (\pi_1, \dots, \pi_T)$ denote the corresponding policy.

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Let $\pi = (\pi_1, \dots, \pi_T)$ denote the corresponding policy.

Then, if the value function \hat{V}_t has \mathfrak{F} -constant K_t , then

► For any history h_t ,

$$|V_t(h_t) - \hat{V}_t(\varphi_t(h_t))|$$

$$\leq \varepsilon_T + \sum_{s=t}^T (\varepsilon_s + K_s \delta_s).$$

► For any history h_t ,

$$|V_t(h_t) - V_t^\pi(h_t)|$$

$$\leq 2[\varepsilon_T + \sum_{s=t}^T (\varepsilon_s + K_s \delta_s)].$$

AIS: Some remarks

In the definition of AIS, we can replace

$$d_{\mathfrak{F}}(\mathbb{P}(\mu_t, \nu_t(\cdot | Z_t = \varphi_t(h_t), U_t = u_t)) \leq \delta_t$$

by

- ▷ $Z_{t+1} = \text{function}(Z_t, Y_t, U_t)$
- ▷ $d_{\mathfrak{F}}(\mathbb{P}(Y_t | H_t = h_t, U_t = u_t), \mathbb{P}(Y_t | Z_t = \varphi_t(h_t), U_t = u_t)) \leq \delta_t.$

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Two ways to interpret the results:

- ▷ Given the information state space \mathcal{Z} , find the best compression $\varphi_t: \mathcal{H}_t \rightarrow \mathcal{Z}$
- ▷ Given any compression function $\varphi_t: \mathcal{H}_t \rightarrow \mathcal{Z}_t$, find the approximation error.

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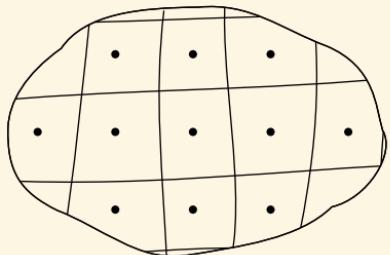
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Results naturally extend to infinite horizon

Some examples

Example 1: Error bounds on state aggregation

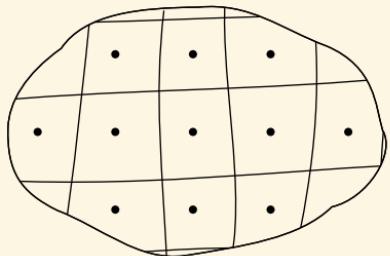


Consider an MDP with state space \mathcal{X} and per-step reward $R_t = r(X_t, U_t)$.

Suppose \mathcal{X} is quantized to a discrete set \mathcal{Z} using $\varphi: \mathcal{X} \rightarrow \mathcal{Z}$.

- ▷ Let $z = \varphi(x)$ denote the label for x .
- ▷ Then $\varphi^{-1}(z)$ denote all states which have label z .

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$\{Z_t\}_{t \geq 1}$ IS AN (ε, δ) AIS

$$\varepsilon = \sup_{(x,u) \in \mathcal{X} \times \mathcal{U}} |r(x, u) - r(\varphi(x), u)|$$

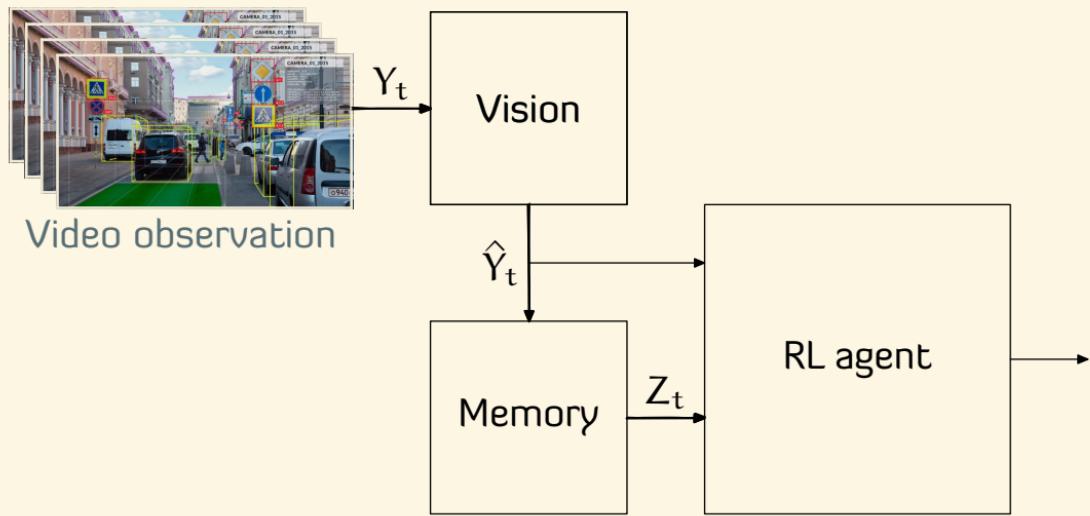
$$\delta = \sup_{(x,u) \in \mathcal{X} \times \mathcal{U}} d_{\mathfrak{F}}(\mathbb{P}(X_+ | X = x, U = u), \mathbb{P}(X_+ | X \in \varphi^{-1}(\varphi(x)), U = u)).$$

or equivalently, $r(\cdot, u)$ has a \mathfrak{F} -constant K_r , $\mathbb{P}(X_+ | X = \cdot, U = u)$ has a \mathfrak{F} -constant K_p , then

$$\varepsilon = K_r D, \quad \delta = K_p D, \quad \text{where } D = \max\{\|x - y\| : \varphi(x) = \varphi(y)\}.$$

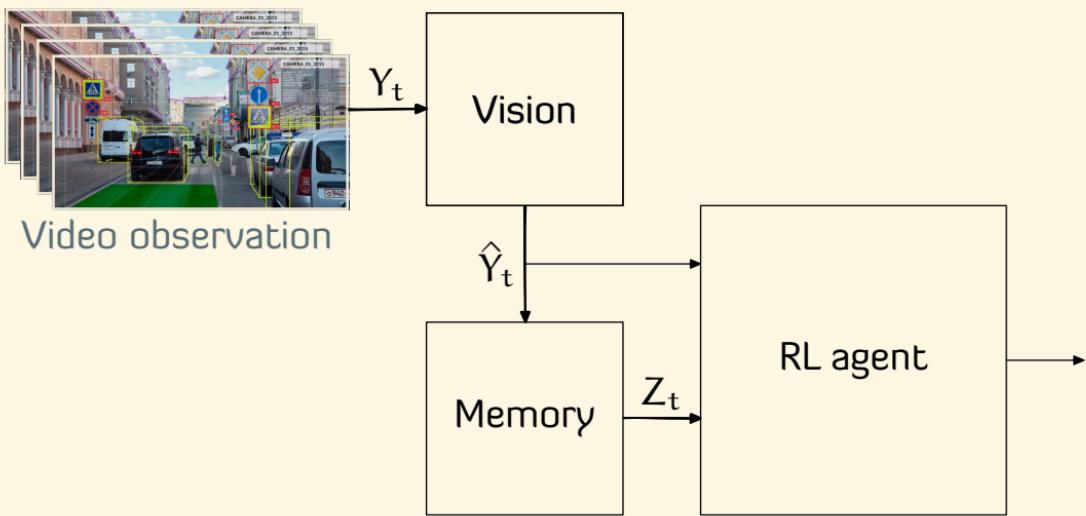
▷ Bertsekas, "Convergence of discretization procedures in dynamic programming," 1975.

Example 2: Approximation bounds for using quantized obs.



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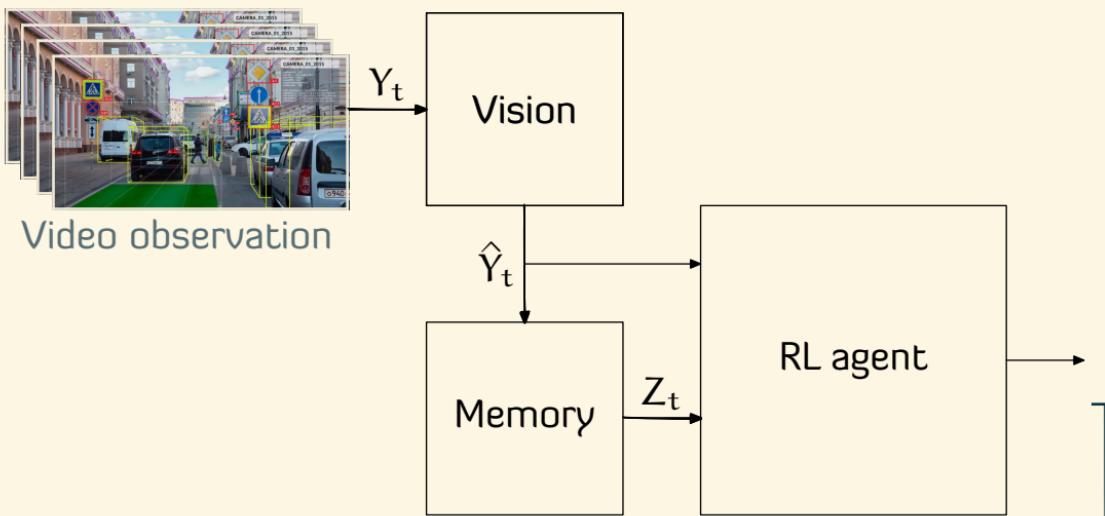
- ▷ Proposed as a heuristic algorithms
- ▷ No performance bounds



▷ Ha, Schmidhuber, "World Models", 2018.

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$$\varepsilon_t = \sup_{h_t, u_t} |\mathbb{E}[R_t | h_t, u_t] - \tilde{R}_t(\varphi_t(h_t), u_t)|$$

$$\delta_t = \sup_{h_t, u_t} d_{\mathfrak{F}}(\mathbb{P}(\hat{Y}_{t+1} | h_t, u_t), \mathbb{P}(\hat{Y}_{t+1} | \varphi_t(h_t), u_t))$$

Example 3: Approximation bounds for mean-field teams

n agents: state X_t^i , control U_t^i .

► Empirical mean-field:

$$M_t(x) = \frac{1}{n} \sum_{i=1}^n \delta_{X_t^i}(x).$$

► Statistical mean-field:

$$\bar{m}_t(x) = \mathbb{P}(X_t^i = x).$$

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► Dynamics

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► Per-step reward

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▷ Info structure: $I_t^i = \{X_t^i, M_t\}$

▷ Infinite population limit: $\tilde{I}_t^i = \{X_t^i, \bar{m}_t\}$.

$$\bar{m}_{t+1} = \mathcal{P}_g m_t,$$

where

$$[\mathcal{P}_g m](y) = \sum_{x,u} m(x) g(u|x) P(y|x, u, m).$$

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(A) $r(x, u, m)$ and $P(y|x, u, m)$ are Lipschitz in x , u , and m .

$\{\bar{m}_t\}_{t \geq 1}$ is an (ε, δ) AIS for expanded info structure, where $\varepsilon, \delta \in O(1/\sqrt{n})$.

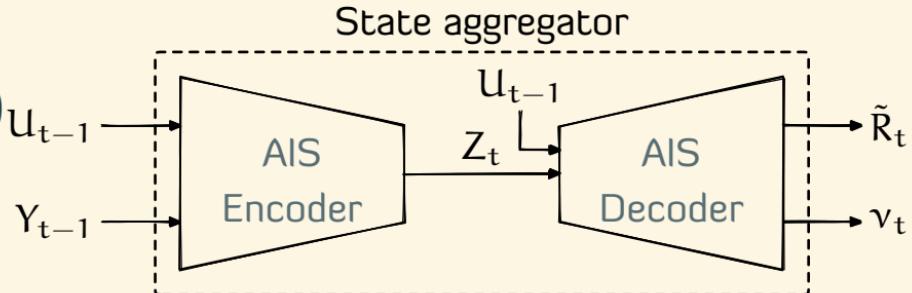
Now to reinforcement learning
for partially observed systems.

Reinforcement learning setup

▷ State aggregator:

$$\mathcal{L}_{AIS} = \alpha_t |\tilde{R}_t - R_t| + (1 - \alpha_t) d_{\mathfrak{F}}(\nu_t, \mu_t)$$

ξ : Parameters of the aggregator
Updated using SGD with LR α_k



Reinforcement learning setup

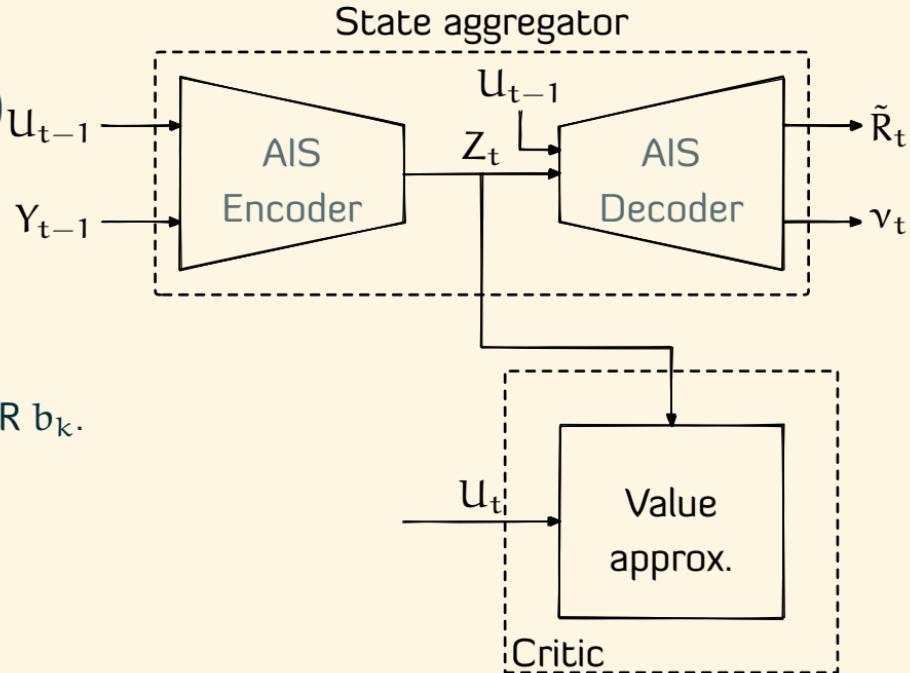
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▷ Value approximator:

$$\varphi$$
: parameters of $Q(z, u)$ approximator.
Updated using TD(0) or TD(λ) with LR b_k .



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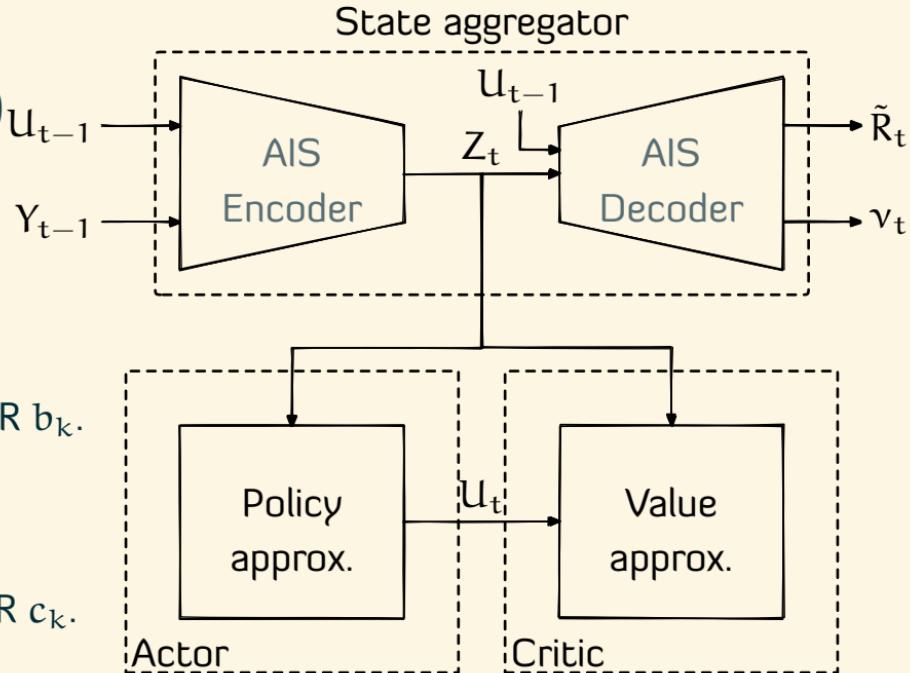
φ : parameters of $Q(z, u)$ approximator.

Updated using TD(0) or TD(λ) with LR b_k .

▷ Policy approximator:

θ : parameters of $\pi(u | z)$

Updated using policy gradient with LR c_k .



Reinforcement learning setup

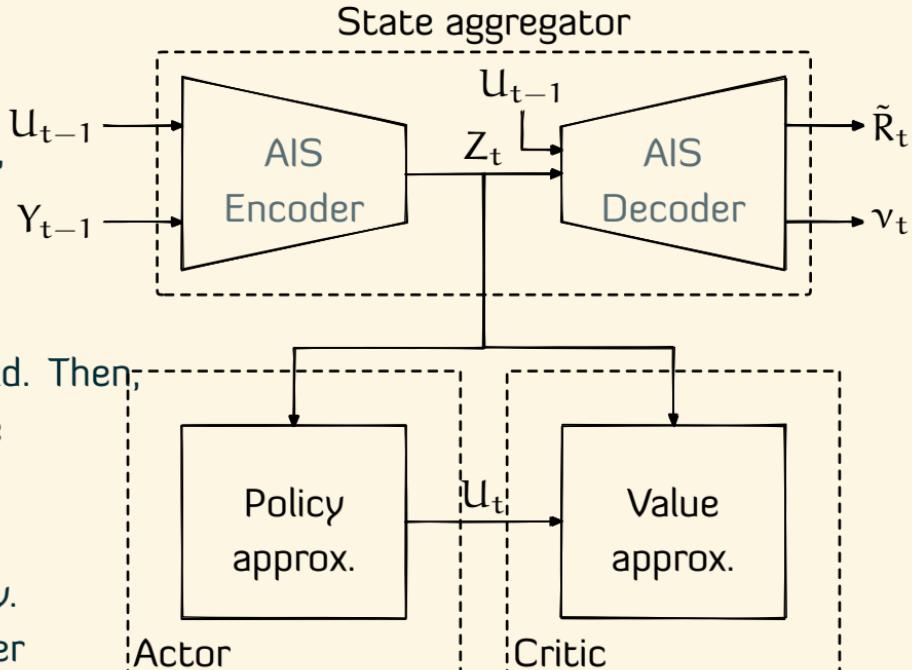
CONVERGENCE RESULT

If the learning rates satisfy conditions for three time-scale stochastic approximation, the compatibility condition

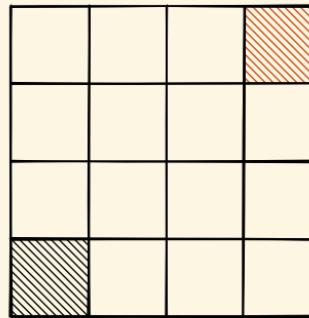
$$\frac{\partial Q(z, u)}{\partial \varphi} = \frac{1}{\pi(u|z)} \frac{\partial \pi(u|z)}{\partial \theta}$$

and additional mild technical conditions hold. Then,

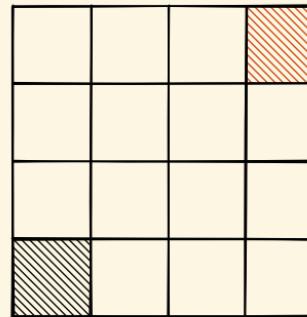
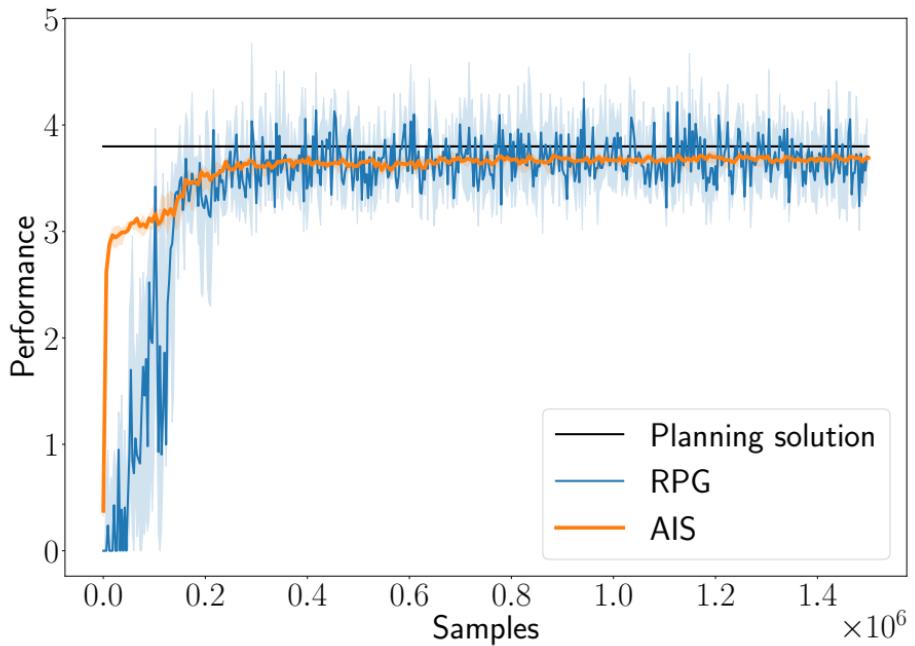
- ▷ State aggregator converges (with some approximation error)
- ▷ The critic converges to the best approximator within the specified family.
- ▷ The actor converges to a local maximizer within the family of policy approximators.



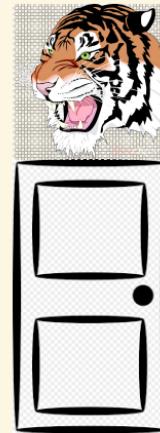
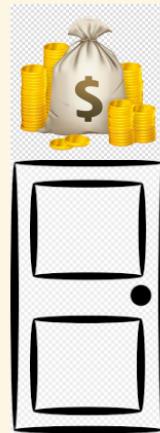
Numerical Results: 4×4 Grid Environment



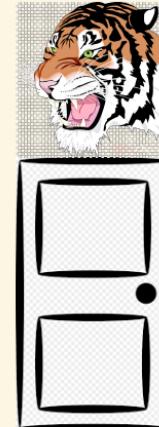
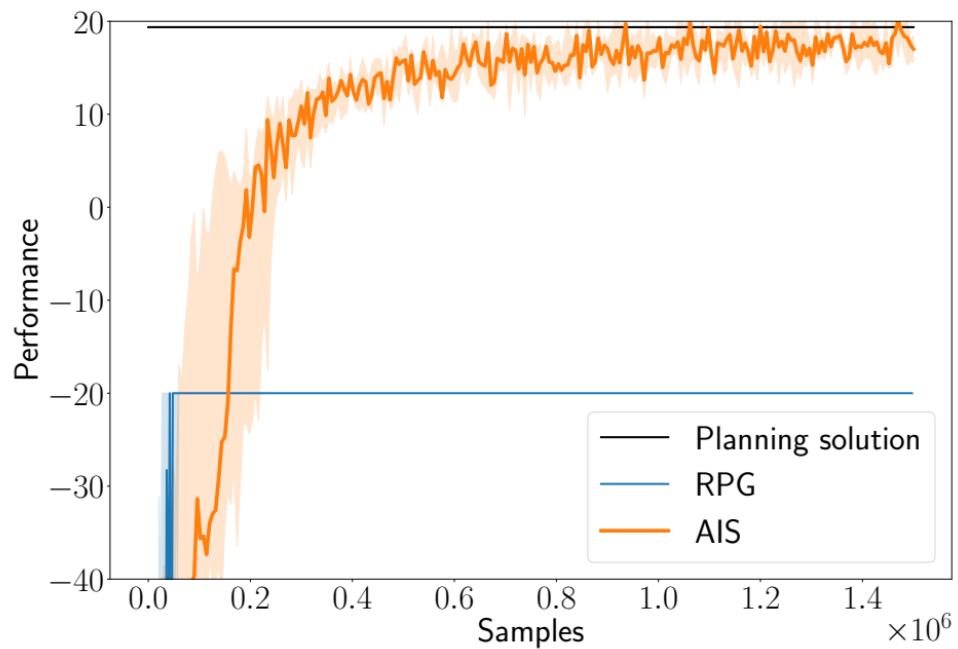
Numerical Results: 4×4 Grid Environment



Numerical Results: Tiger Environment

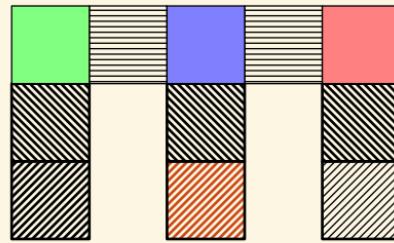


Numerical Results: Tiger Environment

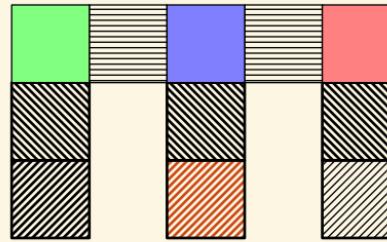
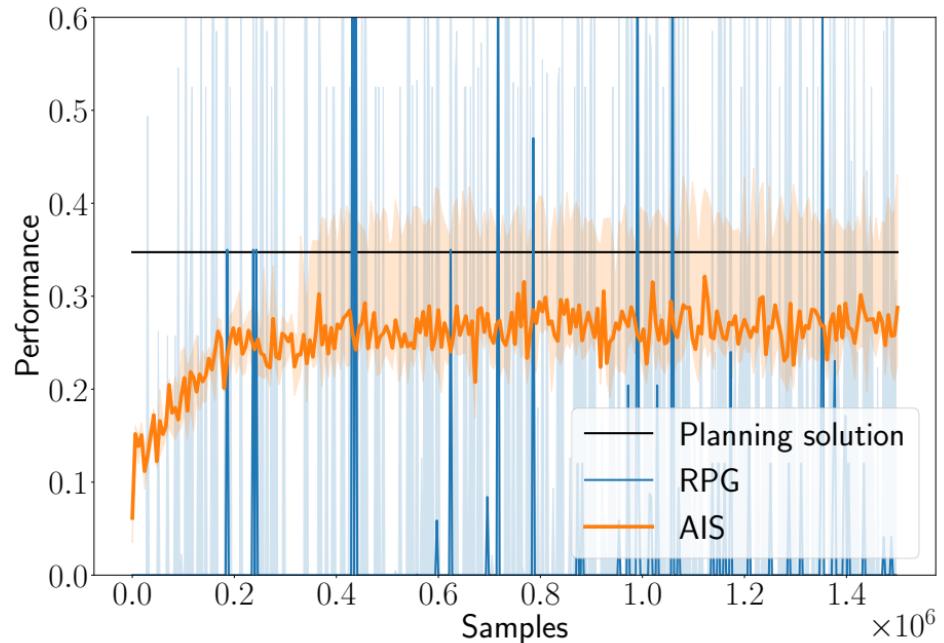


Approx. POMDPs-(Mahajan)

Numerical Results: Cheese Maze Environment

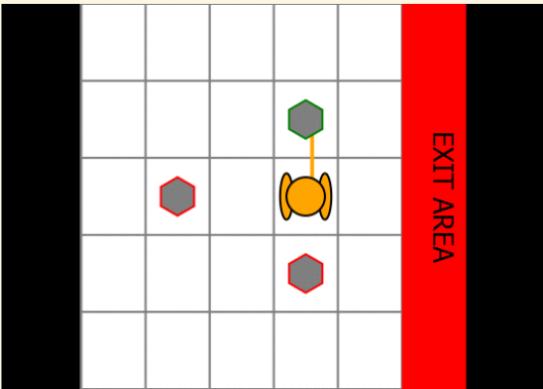


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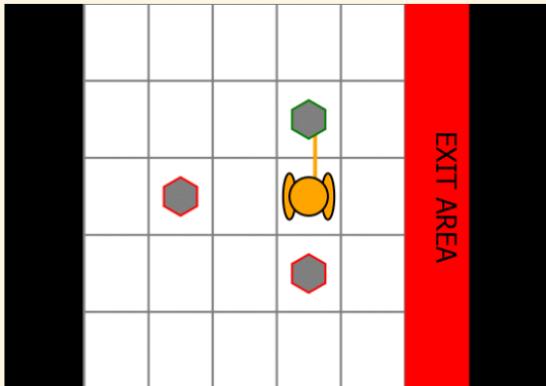
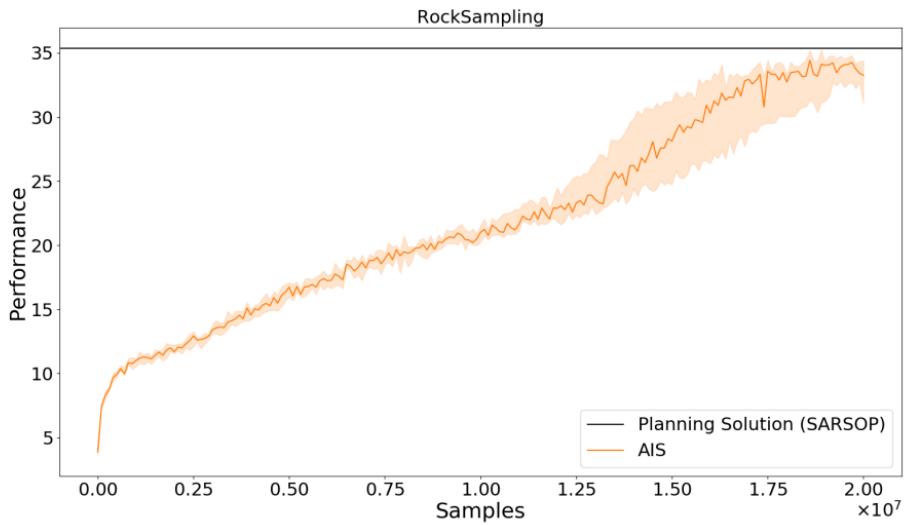


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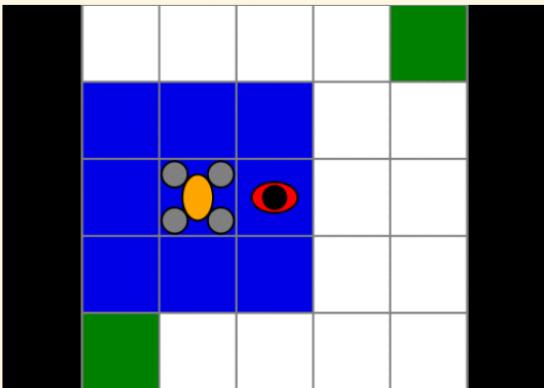
Numerical Results: Rock Sample



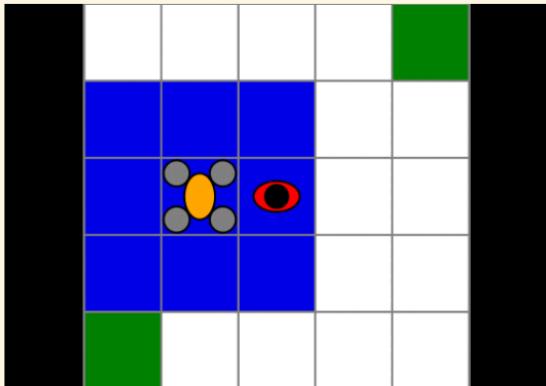
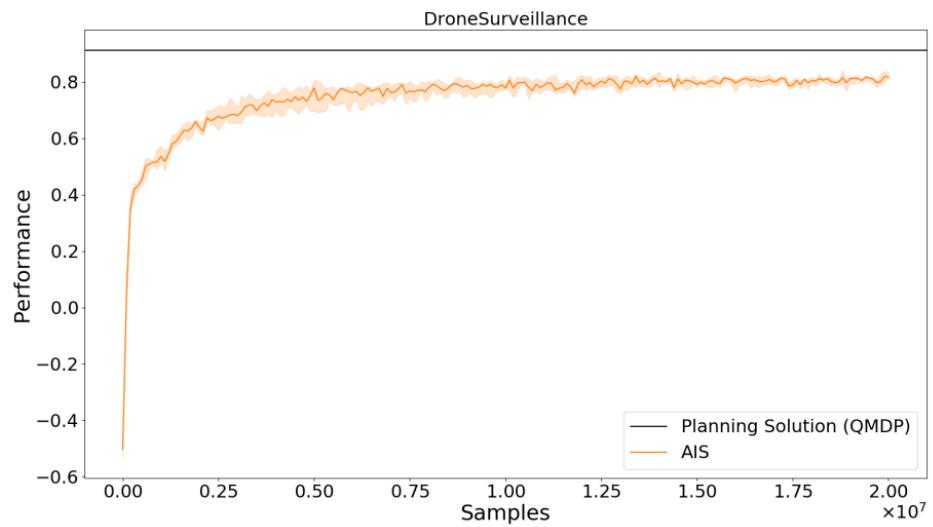
Numerical Results: Rock Sample



Numerical Results: Drone Surveillance



Numerical Results: Drone Surveillance



Summary

Summary

Now let's construct the state space

FORECASTING OUTPUTS IN DISTRIBUTION

$H_t^{(1)} \sim H_t^{(2)}$ if for all future CONTROL inputs $U_{t:T}$,
 $\mathbb{P}(Y_{t:T}^{(1)} | H_t^{(1)}, U_{t:T}) = \mathbb{P}(Y_{t:T}^{(2)} | H_t^{(2)}, U_{t:T})$

Same complexity as identifying the state sufficient for forecasting outputs for the case of perfect observations (which was Step 1 for belief state formulations)

PROPERTIES OF INFORMATION STATE

The info state Z_t at time t is a "compression" of past inputs that satisfies the following:

▷ SUFFICIENT TO PREDICT ITSELF:

$$\mathbb{P}(Z_{t+1} | H_t, U_t) = \mathbb{P}(Z_{t+1} | Z_t, U_t).$$

▷ SUFFICIENT TO PREDICT OUTPUT:

$$\mathbb{P}(Y_t | H_t, U_t) = \mathbb{P}(Y_t | Z_t, U_t).$$

KEY QUESTIONS

- ▷ Can this be used for dynamic programming?
- ▷ What is the right notion of approximations in this framework?

Approx. POMDPs-(Mahajan)



Approx. POMDPs-(Mahajan)

Summary

Now let's construct the state space

Approximate information state

(ε, δ) -APPROXIMATE INFORMATION STATE (AIS)

Given a function class \mathfrak{F} , a compression $\{Z_t\}_{t \geq 1}$ of history (i.e., $Z_t = \varphi_t(H_t)$) is called an $\{(\varepsilon_t, \delta_t)\}_{t \geq 1}$ AIS if there exist:

- ▷ a function $\tilde{R}_t(Z_t, U_t)$, and ▷ a stochastic kernel $v_t(Z_{t+1}|Z_t, U_t)$

such that

- ▷ $|\mathbb{E}[R_t|H_t = h_t, U_t = u_t] - \tilde{R}_t(\varphi_t(h_t), u_t)| \leq \varepsilon_t$

- ▷ For any Borel set A of \mathcal{Z}_{t+1} , define

$$\mu_t(A) = \mathbb{P}(Z_{t+1} \in A | H_t = h_t, U_t = u_t)$$

Then,

$$d_{\mathfrak{F}}(\mu_t, v_t(\cdot | \varphi_t(h_t), u_t)) \leq \delta_t.$$

Summary

Now let's construct the state space

Approximate dynamic programming using AIS

MAIN THEOREM

Given a function class \mathfrak{F} , let $\{Z_t\}_{t \geq 1}$, where $Z_t = \varphi_t(H_t)$, be an $\{(\varepsilon_t, \delta_t)\}_{t \geq 1}$ AIS.

Recursively define the following functions:

$$\hat{V}_{T+1}(z_{T+1}) = 0$$

and for $t \in \{T, \dots, 1\}$:

$$\hat{V}_t(z_t) = \max_{u_t \in \mathcal{U}} \left\{ \tilde{R}_t(z_t, u_t) \right.$$

$$+ \int \hat{V}_{t+1}(z_{t+1}) v_t(dz_{t+1} | z_t, u_t) \Big\}.$$

Let $\pi = (\pi_1, \dots, \pi_T)$ denote the corresponding policy.

Then, if the value function \hat{V}_t has \mathfrak{F} -constant K_t , then

▷ for any history h_t ,

$$\begin{aligned} & |V_t(h_t) - \hat{V}_t(\varphi_t(h_t))| \\ & \leq \varepsilon_T + \sum_{s=t}^T (\varepsilon_s + K_s \delta_s). \end{aligned}$$

▷ for any history h_t ,

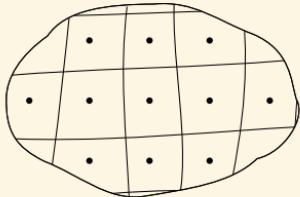
$$\begin{aligned} & |V_t(h_t) - V_t^\pi(h_t)| \\ & \leq 2[\varepsilon_T + \sum_{s=t}^T (\varepsilon_s + K_s \delta_s)]. \end{aligned}$$

Summary

Now let's construct the state space

Approximate dynamic programming using AIS

Example 1: Error bounds on state aggregation



Consider an MDP with state space \mathcal{X} and per-step reward $R_t = r(X_t, U_t)$.

Suppose \mathcal{X} is quantized to a discrete set \mathcal{Z} using $\varphi: \mathcal{X} \rightarrow \mathcal{Z}$.

- ▷ Let $z = \varphi(x)$ denote the label for x .
- ▷ Then $\varphi^{-1}(z)$ denote all states which have label z .

$\{Z_t\}_{t \geq 1}$ IS AN (ε, δ) AIS

$$\varepsilon = \sup_{(x,u) \in \mathcal{X} \times \mathcal{U}} |r(x, u) - r(\varphi(x), u)|$$

$$\delta = \sup_{(x,u) \in \mathcal{X} \times \mathcal{U}} d_{\mathfrak{F}}(\mathbb{P}(X_+ | X = x, U = u), \mathbb{P}(X_+ | X \in \varphi^{-1}(\varphi(x)), U = u)).$$

or equivalently, $r(\cdot, u)$ has a \mathfrak{F} -constant K_r , $\mathbb{P}(X_+ | X = \cdot, U = u)$ has a \mathfrak{F} -constant K_p ,
then

$$\varepsilon = K_r D, \quad \delta = K_p D, \quad \text{where } D = \max\{\|x - y\| : \varphi(x) = \varphi(y)\}.$$

▷ Bertsekas, "Convergence of discretization procedures in dynamic programming," 1975.

Approx. POMDPs-(Mahajan)



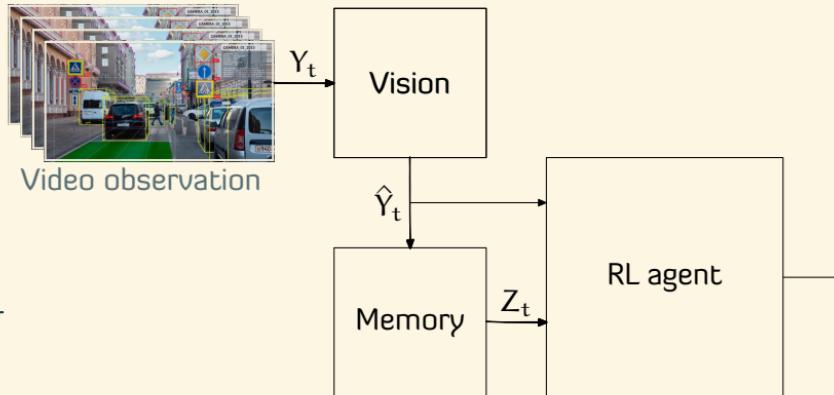
Summary

Now let's construct the state space

Approximate dynamic programming using AIS

Example 2: Approximation bounds for using quantized obs.

- ▷ Proposed as a heuristic algorithms
- ▷ No performance bounds



$\{Z_t\}_{t \geq 1}$ IS AN (ε, δ) AIS

$$\varepsilon_t = \sup_{h_t, u_t} |\mathbb{E}[R_t | h_t, u_t] - \tilde{R}_t(\varphi_t(h_t), u_t)|$$

$$\delta_t = \sup_{h_t, u_t} d_{\mathfrak{F}}(\mathbb{P}(\hat{Y}_{t+1} | h_t, u_t), \mathbb{P}(\hat{Y}_{t+1} | \varphi_t(h_t), u_t))$$

▷ Ha, Schmidhuber, "World Models", 2018.

Approx. POMDPs-(Mahajan)



Summary

Now let's construct the state space.

Approximate dynamic programming using AIS

Example 3: Approximation bounds for mean-field teams

n agents: state X_t^i , control U_t^i .

▷ Dynamics

$$\mathbb{P}(X_{t+1}|X_t, U_t) = \prod_{i=1}^n P(X_{t+1}^i|X_t^i, U_t^i, M_t)$$

▷ Per-step reward

$$R(X_t, U_t) = \frac{1}{n} \sum_{i=1}^n r(X_t^i, U_t^i, M_t)$$

▷ Info structure: $I_t^i = \{X_t^i, M_t\}$

▷ Infinite population limit: $\tilde{I}_t^i = \{X_t^i, \bar{m}_t\}$.

$$\bar{m}_{t+1} = \mathcal{P}_g m_t,$$

where

$$[\mathcal{P}_g m](y) = \sum_{x,u} m(x)g(u|x)\mathbb{P}(y|x, u, m).$$

▷ Empirical mean-field:

$$M_t(x) = \frac{1}{n} \sum_{i=1}^n \delta_{X_t^i}(x).$$

▷ Statistical mean-field:

$$\bar{m}_t(x) = \mathbb{P}(X_t^i = x).$$

Approx. POMDPs-(Mahajan)



(A) $r(x, u, m)$ and $P(y|x, u, m)$ are Lipschitz in x , u , and m .

$\{\bar{m}_t\}_{t \geq 1}$ is an (ε, δ) AIS for expanded info structure, where $\varepsilon, \delta \in \mathcal{O}(1/\sqrt{n})$.

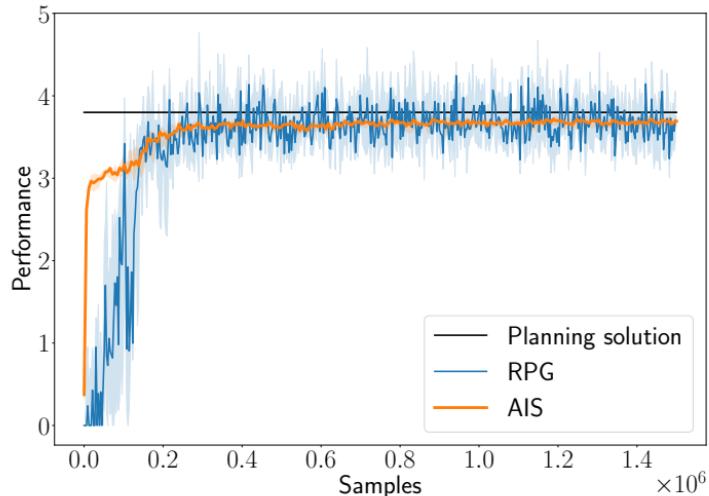
Summary

Now let's construct the state space

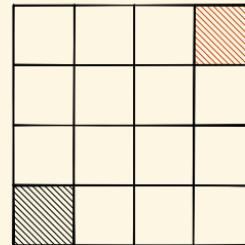
Approximate dynamic programming using AIC

Example 3: Approximation bounds for mean field teams

Numerical Results: 4×4 Grid Environment



Approx. POMDPs-(Mahajan)



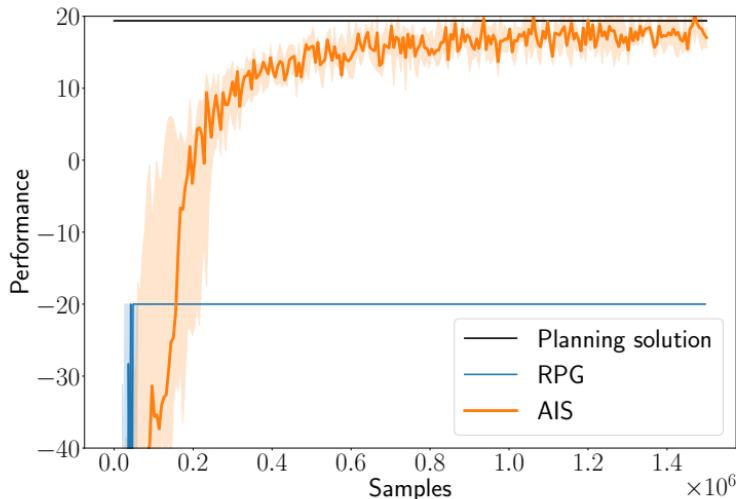
Summary

Now let's construct the state space

Approximate dynamic programming using AIS

Example 3: Approximation bounds for mean field teams

Numerical Results: Tiger Environment



Approx. POMDPs-(Mahajan)



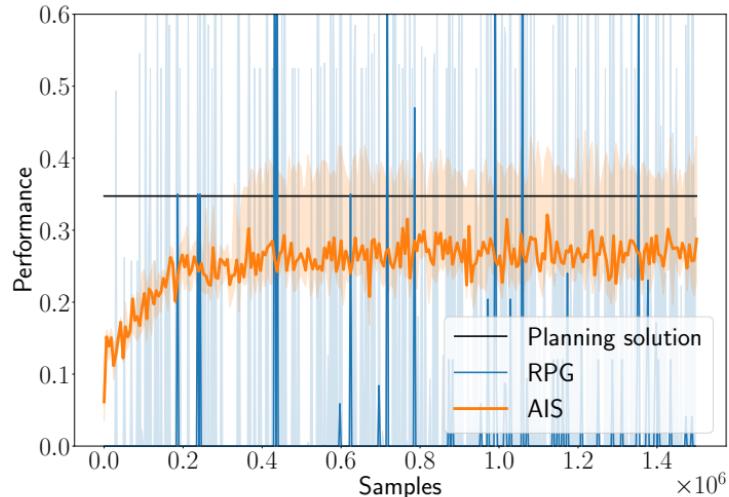
Summary

Now let's construct the state space

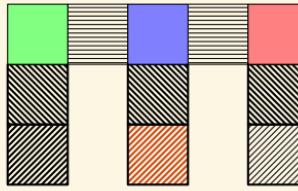
Approximate dynamic programming using AIC

Example 3: Approximation bounds for mean field teams

Numerical Results: Cheese Maze Environment



Approx. POMDPs-(Mahajan)



25

30

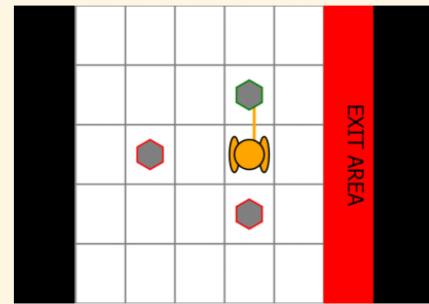
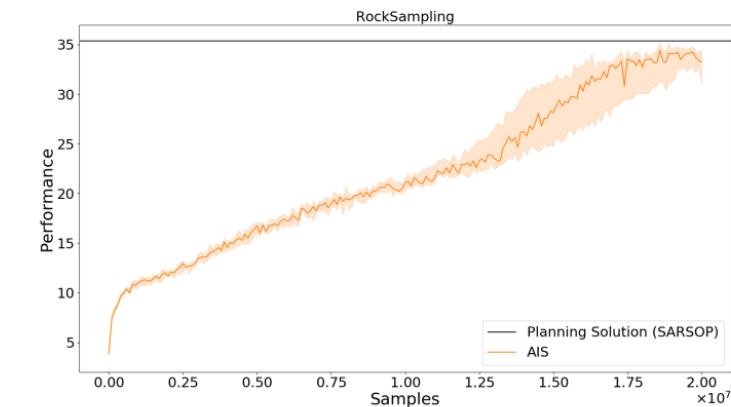
Summary

Now let's construct the state space

Approximate dynamic programming using AIS

Example 3: Approximation bounds for mean field teams

Numerical Results: Rock Sample



Approx. POMDPs-(Mahajan)

27

Approx. POMDPs-(Mahajan)

30

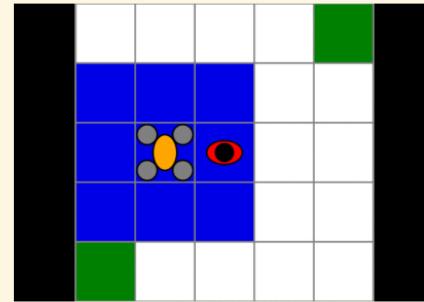
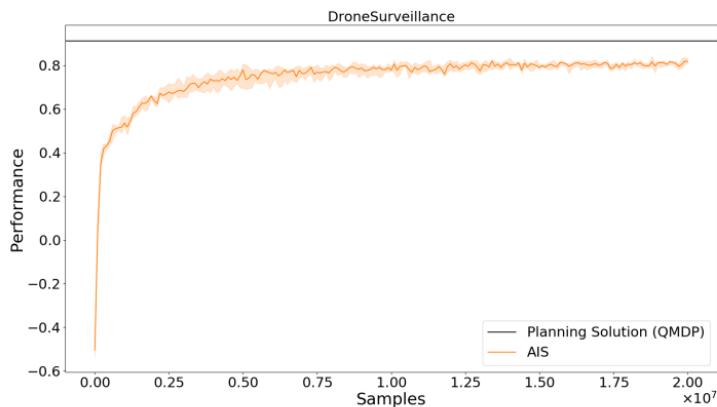
Summary

Now let's construct the state space

Approximate dynamic programming using AIS

Example 3: Approximation bounds for mean field teams

Numerical Results: Drone Surveillance



Approx. POMDPs-(Mahajan)

29

Approx. POMDPs-(Mahajan)

30

Summary

Now let's construct the state space

Approximate dynamic programming using AIS

Example 2: Approximation bounds for mean field teams

Numerical Results: Drone Surveillance

AIS provides a conceptually clean
framework for approximate DP and
online RL in partially observed systems