

Average cost optimal threshold strategies for remote state estimation with communication costs

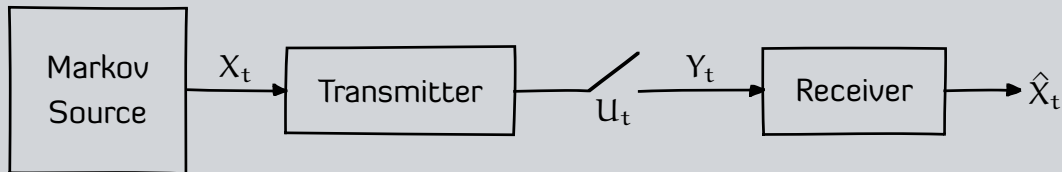
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The communication system



Source ▶ $X_t \in \mathbb{Z}$

▶ Transition matrix P is Toeplitz, i.e., $P_{i,j} = p_{|i-j|}$, where $p_0 \geq p_1 \geq \dots$.

Transmitter $U_t = f_t(X_{1:t}, U_{1:t-1})$ and $Y_t = \begin{cases} X_t & \text{if } U_t = 1 \\ \varepsilon & \text{if } U_t = 0 \end{cases}$

Receiver ▶ $\hat{X}_t = g_t(Y_{1:t})$

▶ Distortion: $d(X_t - \hat{X}_t)$ where $d(e) = d(-e) \leq d(e+1)$

Communication ▶ **Transmission strategy** $f = \{f_t\}_{t=0}^{\infty}$.

Strategies ▶ **Estimation strategy** $g = \{g_t\}_{t=0}^{\infty}$.

The constrained optimization problem

$$\min_{(f,g)} D_\beta(f,g) \quad \text{such that } N_\beta(f,g) \leq \alpha$$

Minimize expected distortion such that expected # of transmissions is less than α

Discounted
setup

$$D_\beta(f,g) = (1 - \beta) \mathbb{E}^{(f,g)} \left[\sum_{t=0}^{\infty} \beta^t d(X_t - \hat{X}_t) \mid X_0 = 0 \right]$$

$$N_\beta(f,g) = (1 - \beta) \mathbb{E}^{(f,g)} \left[\sum_{t=0}^{\infty} \beta^t U_t \mid X_0 = 0 \right]$$

Average cost
setup

$$D_1(f,g) = \limsup_{T \rightarrow \infty} \frac{1}{T} \left[\sum_{t=0}^{T-1} d(X_t - \hat{X}_t) \mid X_0 = 0 \right]$$

$$N_1(f,g) = \limsup_{T \rightarrow \infty} \frac{1}{T} \left[\sum_{t=0}^{T-1} U_t \mid X_0 = 0 \right]$$

Salient Features

- Comparison to Information Theory**
- ▶ As in information theory, the optimization problem may be viewed as minimizing average distortion under an average-power constraint.
 - ▶ Unlike information theory, the source reconstruction must be done in real-time (or with zero delay).
 - ▶ Therefore, classical information theory techniques do not work. Source-channel separation is not optimal.
 - ▶ We use the **decentralized control** approach to real-time communication (following Witsenhausen, Walrand-Varaiya, Teneketzis, . . .)

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 - ▶ Therefore, classical information theory techniques do not work. Source-channel separation is not optimal.
 - ▶ We use the **decentralized control** approach to real-time communication (following Witsenhausen, Walrand-Varaiya, Teneketzis, . . .)

- Comparison to decentralized control**
- ▶ Two decision makers—the transmitter and the receiver.
 - ▶ **(One-sided) nested information structure:**
 - the transmitter knows all the information available to the receiver.
 - ▶ Constrained optimization problem, where the constraint does not depend on the “common information” (i.e., the information at the receiver).

Literature Overview

[Imer-Başar 2005 & 2010]

Fixed number of transmissions for finite horizon LQG setup.

[Lipsa-Martins 2009 & 2011, Molin-Hirche 2009]

Remote estimation with communication cost for finite horizon LQG setup.

[Nayyar-Başar-Teneketzi-Veeravalli 2013]

Remote estimation with communication cost for finite horizon Markov chain setup.
Also considered energy harvesting at the transmitter.

A large literature on [event-driven communication](#) . . .

Assumptions on the model

(A0) $X_t \in \mathbb{Z}$, and $X_0 = 0$.

(A1) The transition matrix is Toeplitz with decaying off-diagonal terms.

$$P = \begin{bmatrix} \ddots & p_0 & \ddots & & & & \\ \cdots & p_1 & p_0 & p_1 & \cdots & & \\ & \ddots & p_1 & p_0 & p_1 & \cdots & \\ & & \ddots & \ddots & p_0 & \ddots & \\ & & & & & & \ddots \end{bmatrix} \quad \text{and} \quad p_0 \geq p_1 \geq p_2 \geq \cdots$$

► Nayyar et al, assumed that the transition matrix was banded, that is, $\exists b$ such that $p_k = 0$, for all $k \geq b$.

(A2) The distortion function is even and increasing on $\mathbb{Z}_{\geq 0}$.

$$\forall e \in \mathbb{Z}_{\geq 0} : \quad d(e) = d(-e) \quad \text{and} \quad d(e) \leq d(e+1).$$

Lagrange Relaxation

$$\min_{(f,g)} D_{\beta}(f, g) \quad \text{such that } N_{\beta}(f, g) \leq \alpha$$

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Lagrange
Relaxation

$$C_{\beta}^*(\lambda) := \inf_{(f,g)} C_{\beta}(f, g; \lambda) \quad \text{where } C_{\beta}(f, g; \lambda) = D_{\beta}(f, g) + \lambda N_{\beta}(f, g)$$

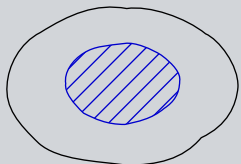
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Search space of
strategies (f, g)

- ▶ **Restrict** the search space of strategies (f, g) by identifying structure of optimal transmission and estimation strategies.
- ▶ **Difficulty**: Non-classical information structure

Structure of optimal estimator (Nayyar et al, 2013)

Transmitted Process Let Z_t denote the most recently transmitted value of the Markov source.

$$Z_0 = 0 \quad \text{and} \quad Z_t = \begin{cases} X_t & \text{if } U_t = 1; \\ Z_{t-1} & \text{if } U_t = 0. \end{cases}$$

The estimator can keep track of Z_t as follows:

$$Z_0 = 0 \quad \text{and} \quad Z_t = \begin{cases} Y_t & \text{if } Y_t \neq \varepsilon; \\ Z_{t-1} & \text{if } Y_t = \varepsilon. \end{cases}$$

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Theorem 1 The process $\{Z_t\}_{t=0}^{\infty}$ is a sufficient statistic at the estimator and an optimal estimation strategy is given by

$$\hat{X}_t = g_t^*(Z_t) = Z_t \quad (*)$$

Remark ▶ The optimal estimation strategy is **time-homogeneous** and can be specified in closed form.

Structure of optimal transmitter (Nayyar et al, 2013)

Error process Let $E_t = X_t - Z_{t-1}$ denote the error process. $\{E_t\}_{t=0}^{\infty}$ is a controlled Markov process where

$$E_0 = 0 \quad \text{and} \quad \mathbb{P}(E_{t+1} = n \mid E_t = e, U_t = u) = \begin{cases} P_{0n}, & \text{if } u = 1; \\ P_{en}, & \text{if } u = 0. \end{cases}$$

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Theorem 2 When the estimation strategy is of the form (\star) , then $\{E_t\}_{t=0}^{\infty}$ is a sufficient statistic at the transmitter.

Furthermore, an optimal transmission strategy is characterized by a time-varying threshold $\{k_t\}_{t=0}^{\infty}$, i.e.,

$$U_t = f_t(E_t) = \begin{cases} 1 & \text{if } |E_t| \geq k_t; \\ 0 & \text{if } |E_t| < k_t. \end{cases}$$

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Proof idea ▶ The proof of [Nayyar et al, 2013] was based on some majorization inequalities of [Hajek et al, 2009] for distributions with finite support.
▶ We extend these inequalities to distributions over integers using results of [Wang-Woo-Madiman, 2014].

Infinite horizon setup (for Lagrange relaxation)

- Main idea** ▶ Based on Thm 1, restrict attention to **time-homogeneous** estimation strategy

$$\hat{X}_t = g_t^*(Z_t) = Z_t$$

- ▶ Consider the problem of finding the “best response” estimation strategy.

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- ▶ Centralized stochastic control problem with countable state space and unbounded cost.
- ▶ Standard MDP results apply under mild technical assumptions.

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- Assum (A₃)** For every $\lambda \geq 0$, there exists a function $w : \mathbb{Z} \rightarrow \mathbb{R}$ and positive and finite constants μ_1 and μ_2 such that for all $e \in \mathbb{Z}$, we have that

$$\max\{\lambda, d(e)\} \leq \mu_1 w(e)$$

$$\max \left\{ \sum_{n=-\infty}^{\infty} P_{en} w(n), \sum_{n=-\infty}^{\infty} P_{0n} w(n) \right\} \leq \mu_2 w(e).$$

Structure of optimal transmitter for infinite horizon

Structure Under assumption (A3), optimal transmission strategy is characterized by **time-homogeneous threshold k** , i.e.,

$$u_t = f(E_t) = \begin{cases} 1 & \text{if } |E_t| \geq k; \\ 0 & \text{if } |E_t| < k. \end{cases}$$

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Dynamic program For $\beta \in (0, 1)$, the optimal strategy is determined by the unique fixed point of the following DP:

$$V_\beta(e; \lambda) = \min \left\{ \begin{array}{l} (1 - \beta)\lambda + \beta \sum_{n=-\infty}^{\infty} P_{0n} V_\beta(n; \lambda), \quad \text{Transmit} \\ (1 - \beta)d(e) + \beta \sum_{n=-\infty}^{\infty} P_{en} V_\beta(n; \lambda) \end{array} \right\} \quad \text{Don't Transmit}$$

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Lagrange relaxation Let $f_\beta^*(\cdot; \lambda)$ be the time-homogeneous optimal transmission strategy.

$$C_\beta^*(\lambda) := \inf_{(f, g)} C_\beta(f, g; \lambda) = C_\beta(f_\beta^*, g^*; \lambda) = V_\beta(0; \lambda)$$

The SEN Conditions and the long-term average setup

- SEN Conditions** For any $\lambda \geq 0$, the value function $V_\beta(\cdot; \lambda)$ satisfy the SEN condition:
- (S1) There exists a reference state $e_0 \in \mathbb{Z}$ such that $V_\beta(e_0; \lambda) < \infty$ for all $\beta \in (0, 1)$.
 - (S2) Define $h_\beta(e; \lambda) = (1 - \beta)^{-1}[V_\beta(e; \lambda) - V_\beta(e_0; \lambda)]$. There exists a function $K_\lambda : \mathbb{Z} \rightarrow \mathbb{R}$ such that $h_\beta(e; \lambda) \leq K_\lambda(e)$ for all $e \in \mathbb{Z}$ and $\beta \in (0, 1)$.
 - (S3) There exists a non-negative (finite) constant L_λ such that $-L_\lambda \leq h_\beta(e; \lambda)$ for all $e \in \mathbb{Z}$ and $\beta \in (0, 1)$.

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Vanishing discount approach Let $f_1^*(\cdot; \lambda)$ be any limit point of $f_\beta^*(\cdot; \lambda)$ as $\beta \uparrow 1$. Then the time-homogeneous transmission strategy $f_1^*(\cdot; \lambda)$ is optimal for $\beta = 1$ (the long-term average setup).

Furthermore, the performance of this optimal strategy is

$$C_1^*(\lambda) := \inf_{(f, g)} C_1(f, g; \lambda) = C_1(f_1^*, g^*; \lambda) = \lim_{\beta \uparrow 1} V_\beta(0; \lambda) = \lim_{\beta \uparrow 1} C_\beta^*(\lambda).$$

Performance of a threshold based strategy

Threshold-based strategy

We analyze the performance of $(f^{(k)}, g^*)$, where

$$f^{(k)}(e) := \begin{cases} 1, & \text{if } |e| \geq k; \\ 0, & \text{if } |e| < k. \end{cases}$$

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Cost until first transmission Define $S^{(k)} = \{e \in \mathbb{Z} : |e| \leq k - 1\}$ and let $\tau^{(k)}$ be the stopping time when the Markov process starting at state 0 at time $t = 0$ escapes the set $S^{(k)}$.

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$$\text{Define } L_{\beta}^{(k)} := \mathbb{E} \left[\sum_{t=0}^{\tau^{(k)}-1} \beta^t d(E_t) \middle| E_0 = 0 \right]$$

$$M_{\beta}^{(k)} := \frac{1 - \mathbb{E}[\beta^{\tau^{(k)}} | E_0 = 0]}{1 - \beta}$$

and

$$L_1^{(k)} := \mathbb{E} \left[\sum_{t=0}^{\tau^{(k)}-1} d(E_t) \middle| E_0 = 0 \right]$$

$$M_1^{(k)} := \mathbb{E}[\tau^{(k)} - 1 | E_0 = 0]$$

Performance of a threshold based strategy (cont.)

Renewal
relationships

$$D_{\beta}^{(k)} := D_{\beta}(f^{(k)}, g^*) = \frac{L_{\beta}^{(k)}}{M_{\beta}^{(k)}}$$

$$N_{\beta}^{(k)} := N_{\beta}(f^{(k)}, g^*) = \frac{1}{M_{\beta}^{(k)}} - (1 - \beta)$$

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Vanishing
discount
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$$L_1^{(k)} = \lim_{\beta \uparrow 1} L_{\beta}^{(k)}, \quad M_1^{(k)} = \lim_{\beta \uparrow 1} M_{\beta}^{(k)}.$$

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$$D_1^{(k)} = \lim_{\beta \uparrow 1} D_{\beta}^{(k)} = \frac{L_1^{(k)}}{M_1^{(k)}}$$

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Performance of a threshold based strategy: Computations

Analytic
expressions
for performance

Let $P^{(k)}$ and $Q_{\beta}^{(k)}$ be square matrices and $d^{(k)}$ is a column vector indexed by $S^{(k)}$ defined as follows:

$$P_{ij}^{(k)} := P_{ij}, \quad \forall i, j \in S^{(k)},$$

$$Q_{\beta}^{(k)} := [I_{2k-1} - \beta P^{(k)}]^{-1},$$

$$d^{(k)} := [d(-k+1), \dots, d(k-1)]^T$$

Then,

$$L_{\beta}^{(k)} = [Q_{\beta}^{(k)}]_0 d^{(k)} \quad \text{and} \quad M_{\beta}^{(k)} = [Q_{\beta}^{(k)}]_0 \mathbf{1}_{2k-1}.$$

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$D_{\beta}^{(k)}$ and $N_{\beta}^{(k)}$ can be computed using these expressions.

Optimal strategy for the Lagrange relaxation

Some inequalities

$$L_{\beta}^{(k)} < L_{\beta}^{(k+1)}, \quad M_{\beta}^{(k)} < M_{\beta}^{(k+1)}, \quad D_{\beta}^{(k)} < D_{\beta}^{(k+1)}.$$

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$$C_{\beta}^{(k)}(\lambda) := C(f^{(k)}, g^*; \lambda) = D_{\beta}^{(k)} + \lambda N_{\beta}^{(k)}$$



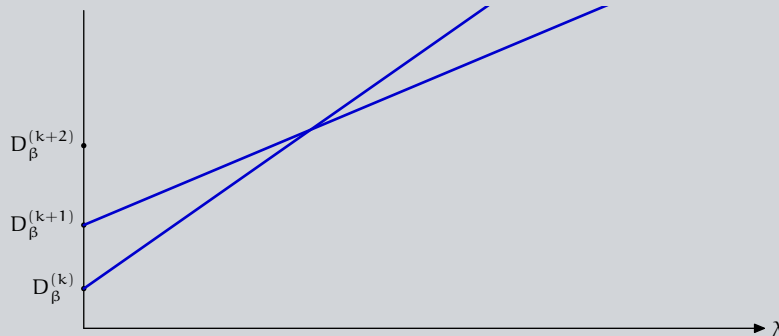
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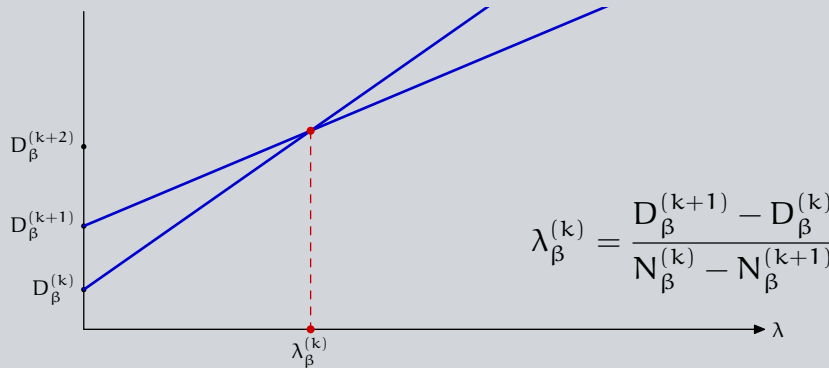
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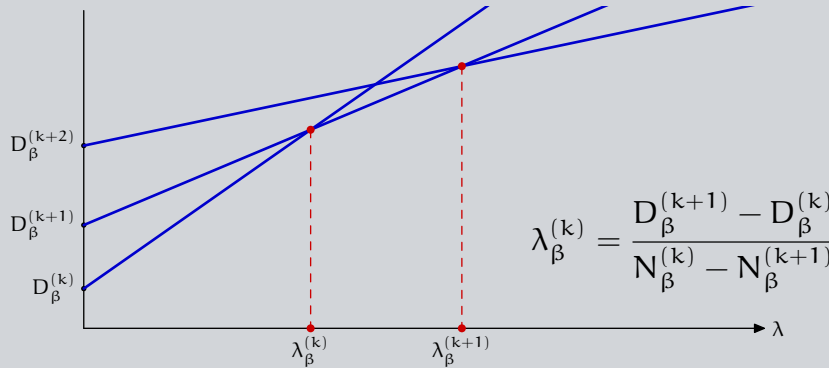
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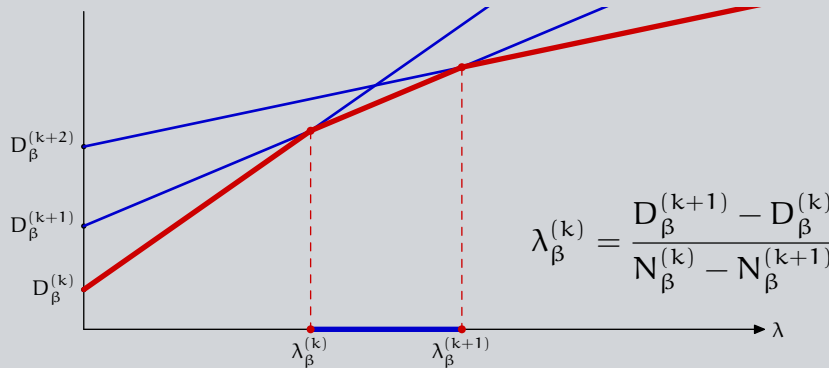
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Optimal performance

- ▶ For all $\lambda \in (\lambda_{\beta}^{(k)}, \lambda_{\beta}^{(k+1)}]$ the threshold strategy $f^{(k+1)}$ is optimal.
- ▶ $C_{\beta}^*(\lambda) = \min_{k \in \mathbb{Z}} C_{\beta}^{(k)}$ is piecewise linear, continuous, concave, and increasing function of λ .

Back to the constrained optimization problem

- Bernoulli randomized strategy** Let $\theta \in [0, 1]$ and f_1 and f_2 be two stationary strategies. The **Bernoulli randomized strategy** (f_1, f_2, θ) randomizes between f_1 and f_2 at each stage, choosing f_1 with probability θ and f_2 with probability $(1 - \theta)$.
- Simple rand. strategy** A Bernoulli randomized strategy (f_1, f_2, θ) is **simple** if the actions prescribed by f_1 and f_2 differ only at one state.

Main result

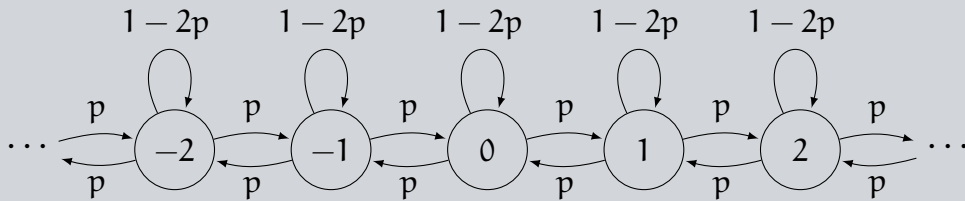
Define $k_\beta^* = \sup\{k \in \mathbb{Z}_{\geq 0} : N_\beta^{(k)} \geq \alpha\}$ and let θ be such that

$$\theta N_\beta^{(k_\beta^*)} + (1 - \theta) N_\beta^{(k_\beta^* + 1)} = \alpha$$

Then, the Bernoulli simple randomized strategy $(f^{(k_\beta^*)}, f^{(k_\beta^* + 1)}, \theta)$ is optimal for the constrained optimization problem for $\beta \in (0, 1]$.

An example: Symmetric birth-death Markov Chain

$$P_{ij} = \begin{cases} p, & \text{if } |i - j| = 1; \\ 1 - 2p, & \text{if } i = j; \\ 0, & \text{otherwise,} \end{cases} \quad \text{where } p \in (0, \frac{1}{2}), \quad d(e) = |e|$$



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Discounted cost Let $K_\beta = -2 - (1 - \beta)/\beta p$ and $m_\beta = \cosh^{-1}(-K_\beta/2)$.

$$D_\beta^{(k)} = \frac{\sinh(km_\beta) - k \sinh(m_\beta)}{2 \sinh^2(km_\beta/2) \sinh(m_\beta)}$$

$$N_\beta^{(k)} = \frac{2\beta p \sinh^2(m_\beta/2) \cosh(km_\beta)}{\sinh^2(km_\beta/2)} - (1 - \beta)$$

Average cost $D_1^{(k)} = \frac{k^2 - 1}{3k}$ and $N_1^{(k)} = \frac{2p}{k^2}$

An example: Symmetric birth-death Markov Chain

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Discounted cost Let $K_\beta = -2 - (1 - \beta)/\beta p$ and $m_\beta = \cosh^{-1}(-K_\beta/2)$.

$$D_\beta^{(k)} = \frac{\sinh(km_\beta) - k \sinh(m_\beta)}{2 \sinh^2(km_\beta/2) \sinh(m_\beta)}$$

$$N_\beta^{(k)} = \frac{2\beta p \sinh^2(m_\beta/2) \cosh(km_\beta)}{\sinh^2(km_\beta/2)} - (1 - \beta)$$

$\lambda_\beta^{(k)}$ can be computed in terms of $D_\beta^{(k)}$ and $N_\beta^{(k)}$.

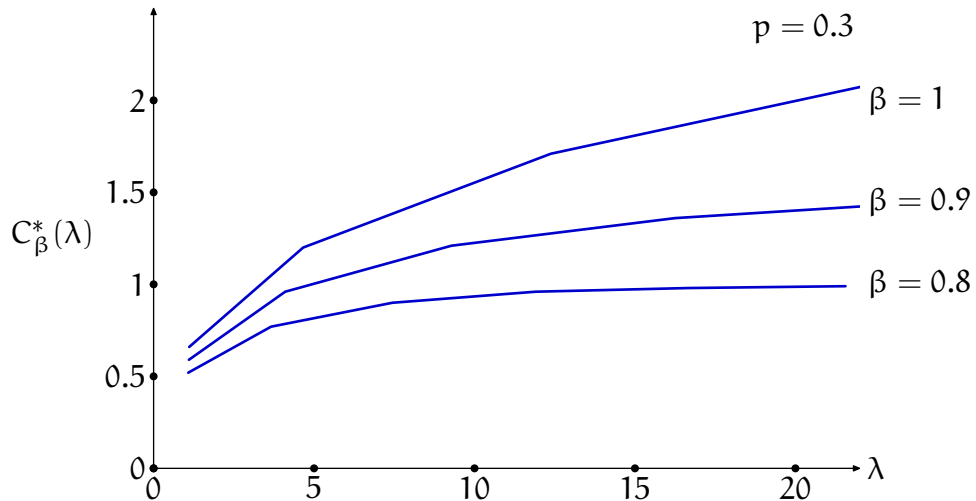
Average cost $D_1^{(k)} = \frac{k^2 - 1}{3k}$ and $N_1^{(k)} = \frac{2p}{k^2}$

$$\lambda_1^{(k)} = \frac{k(k+1)(k^2 + k + 1)}{6p(2k + 1)}$$

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$$k_\beta^* = \sup \left\{ k \in \mathbb{Z}_{\geq 0} : \frac{\sinh^2(m_\beta/2) \cosh(km_\beta)}{\sinh^2(km_\beta/2)} \geq \frac{1 + \alpha - \beta}{2\beta p} \right\}$$

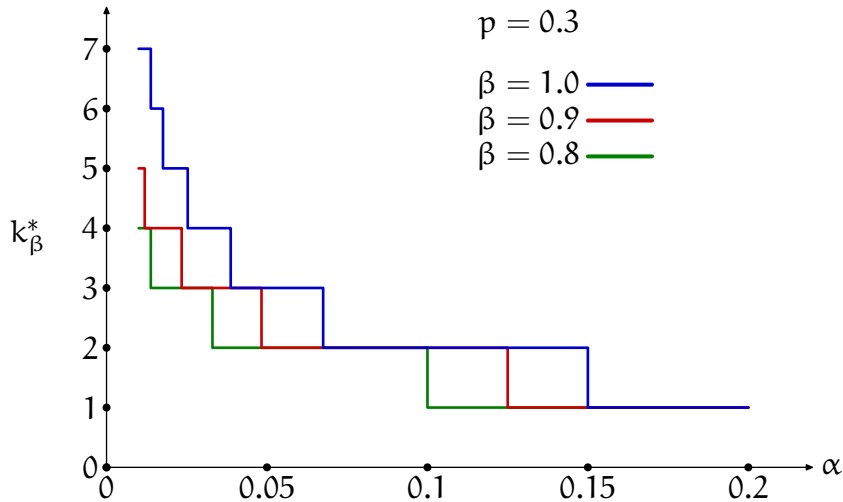
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$$k_1^* = \left\lfloor \sqrt{\frac{2p}{\alpha}} \right\rfloor$$

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Summary and Conclusion

- Problem formulation**
- ▶ **Real-time** transmission of a Markov source under constraints on the number of transmissions.
 - ▶ Investigated both discounted and average cost infinite horizon setups.
 - ▶ Modeled as a **decentralized stochastic control** problem with two decision maker.
 - ▶ As long as the transmitter uses a symmetric threshold based strategy, the estimation strategy does not depend on the transmission strategy.
 - ▶ The problem of find the “best response” transmitter is a centralized stochastic control problem.
- Main results**
- ▶ Simple Bernoulli randomized strategies $(f^{(k^*)}, f^{(k^*+1)}, \theta)$ are optimal.
 - ▶ k^* and θ can be computed easily.