

On Learning Whittle Index Policy for Restless Bandits With Scalable Regret

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Abstract—Reinforcement learning is an attractive approach to learn good resource allocation and scheduling policies based on data when the system model is unknown. However, the cumulative regret of most reinforcement learning (RL) algorithms scales as $\tilde{O}(S\sqrt{AT})$, where S is the size of the state space, A is the size of the action space, T is the horizon, and the $\tilde{O}(\cdot)$ notation hides logarithmic terms. Due to the linear dependence on the size of the state space, these regret bounds are prohibitively large for resource allocation and scheduling problems. In this article, we present a model-based RL algorithm for such problems which has scalable regret. In particular, we consider a restless bandit model, and propose a Thompson-sampling-based learning algorithm which is tuned to the underlying structure of the model. We present two characterizations of the regret of the proposed algorithm with respect to the Whittle index policy. First, we show that for a restless bandit with n arms and at most m activations at each time, the regret scales either as $\tilde{O}(mn\sqrt{T})$ or $\tilde{O}(n^2\sqrt{T})$ depending on the reward model. Second, under an additional technical assumption, we show that the regret scales as $\tilde{O}(n^{1.5}\sqrt{T})$ or $\tilde{O}(\max\{m\sqrt{n}, n\}\sqrt{T})$. We present numerical examples to illustrate the salient features of the algorithm.

Index Terms—Reinforcement learning (RL), restless bandits (RBs), Thompson sampling, Whittle index.

I. INTRODUCTION

RESOURCE allocation and scheduling problems arise in control of networked systems. Examples include opportunistic scheduling in networks [1], [2], [3]; link scheduling in machine-type communication [4]; user allocation in mmWave networks [5]; channel allocation in networks [6]; source selection in peer-to-peer networks [7]; opportunistic spectrum access [8], [9], [10]; demand response in smart grids [11], [12]; dynamic routing in multi-UAVs [13]; operator allocation in multirobot systems [14], etc.

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Due to the curse of dimensionality, finding an optimal solution in such resource allocation and scheduling problems is computationally prohibitive [15]. Restless bandits (RBs) [16] have emerged as a popular solution heuristic for such problems. The RB framework is motivated by the *rested* multiarmed bandit problem considered in the seminal work of Gittins [17], who showed that the optimal strategy for the rested multiarmed bandit problem is of the *index type*: one can compute an index for each state of each alternative (also called an arm), and choose the alternative with the highest index. In general, such index-type policies are not optimal for RBs. In fact, computing the optimal policy for RBs is PSPACE hard [15]. However, as argued in [16], an index-type policy (now known as the Whittle index) can be a useful heuristic if a technical condition known as *indexability* is satisfied. The Whittle index policy is optimal for some specific models [6], [17], [18]. There is also strong empirical evidence to suggest that the Whittle index policy performs close to optimal in various settings [19], [20], [21], [22], [23], [24]. For these reasons, the RB framework has been applied in a variety of resource allocation and scheduling applications referenced above.

In all the above references, it is assumed that the system model is known perfectly. In many real-world applications, there is often uncertainty about the system model. In such situations, reinforcement learning (RL) is an attractive alternative. In recent years, there have been many papers which investigated RL for RBs [25], [26], [27], [28], [29]. Most of these learn the Q-function associated with the average reward/cost optimality equation parameterized by the activation cost λ and use it to asymptotically learn the Whittle index.

A common measure of performance of an RL algorithm is *regret*, which measures the difference in performance of the learning algorithm that does not *a priori* know the system model with the performance of a baseline policy that knows the model. However, the regret is not characterized in existing literature on RL for RBs [25], [26], [27], [28], [29].

There are some results on characterizing regret for some specific instances of RBs: the special case of rested bandits [30], a model of multiclass queues arising in mobile edge computing [31], and a model for scheduling when to observe uncontrolled Markov chains arising in opportunistic spectrum access in cognitive radios [32], [33], [34], [35], [36], [37]. The main contribution of this article is to characterize the regret of a general RL algorithm for general RBs.

It is not possible to directly use existing RL algorithms that achieve near optimal regret in RBs. To explain why this is the case, we provide a short overview of characterizing regret in

TABLE I

COMPARISON OF THE REGRET BOUNDS OF VARIOUS ALGORITHMS

Algorithm	Algorithm Type	Regret ^a	Regret Type
UCRL2 [38]	OUU	$\tilde{O}(DS\sqrt{AT})$	Frequentist
REGAL [39]	OUU	$\tilde{O}(HS\sqrt{AT})$	Frequentist
SCAL [40]	OUU	$\tilde{O}(H\sqrt{\Gamma SAT})$	Frequentist
[41]	TS	$\tilde{O}(D\sqrt{SAT})^b$	Frequentist
[42]	TS	$\tilde{O}(\sqrt{HSAT})$	Frequentist
TSDE [43]	TS	$\tilde{O}(HS\sqrt{AT})$	Bayesian

^a In the column on regret bounds, Γ is the maximum number of states that can be reached from any state, D is the diameter of the MDP, H is the span of the bias of the MDP. These are related as $\Gamma \leq S$ and $H \leq D$ (established in [39]).

^b It is pointed out in [42] that there is a mistake in the proof in [41] and it is suggested that the bound of [41] may be loose by a factor of \sqrt{S} .

RL. Consider a general Markov decision process (MDP) with finite state space of size S and finite action space of size A . It is shown in [38] that no learning algorithm can achieve a regret of less than $\tilde{\Omega}(\sqrt{SADT})$, where D is the diameter of the underlying MDP and T is the time horizon for which the system runs. Several classes of algorithms have been proposed in literature which achieve this lower bound up to a factor of \sqrt{S} and logarithmic terms. Broadly speaking, these regret optimal RL algorithms fall into two classes: 1) optimism under uncertainty (OUU) and 2) Thompson sampling (TS). Two types of regret bounds are provided: 1) frequentist regret, which is a bound on the worst-case regret with high probability and 2) Bayesian regret, which is a bound on the average regret (with respect to a prespecified prior). A summary of the regret bounds for various algorithms is shown in Table I.

Each of these state-of-the-art algorithms has a regret that scales approximately as $\tilde{O}(S\sqrt{AT})$, which is prohibitively large when translated to the RB setting for reasons explained below. Consider an RB with n arms where at most m arms can be activated at a time. Let S^i denote the size of the state space of arm $i \in \{1, \dots, n\}$. Such an RB can be modeled as an MDP, where the size of the state space is $\prod_{i=1}^n S^i$ and the size of the action space is $\binom{n}{m}$. Thus, the regret of using any of the algorithms described in Table I on RBs will be $\tilde{O}\left(\prod_{i=1}^n S^i \sqrt{\binom{n}{m} T}\right)$, which grows exponentially with the number n of arms. In this article, we provide a more nuanced characterization of the scaling of the regret with the number of alternatives.

In particular, we propose a Thompson-sampling-based learning algorithm for RB, which we call as RB-TSDE. This algorithm is inspired from the Thompson sampling with dynamic episodes (TSDE) algorithm [43]. We show that for an RB with n arms, where m of them can be chosen at a time, RB-TSDE has a Bayesian regret (with respect to the Whittle index policy with known dynamics) of $\tilde{O}(n^2\sqrt{T})$ or $\tilde{O}(nm\sqrt{T})$ depending on the assumptions on the per-step reward. Under an additional technical assumption, we obtain an alternative regret bound of $\tilde{O}(n^{1.5}\sqrt{T})$ or $\tilde{O}(\max\{m\sqrt{n}, n\}\sqrt{T})$.

The rest of this article is organized as follows. In Section II, we formulate the learning problem for RB when the state transition probabilities of all arms are unknown and present the main results. In Section III, we present the TSDEs for RB

and provide an upper bound on the regret. In Section IV, we provide the proof outline and defer the details to the Appendix. In Section V, we demonstrate a numerical example of the regret of our algorithm. In Section VI, we discuss relaxation and sufficient conditions of some of the assumptions, in addition to comparison with the optimal policy. Finally, Section VII concludes this article.

Notation: We use uppercase variables S, A , etc., to denote random variables, the corresponding lower variables (s, a , etc.) to denote their realizations, and corresponding calligraphic letters (\mathcal{S}, \mathcal{A} , etc.) to denote set of realizations. Subscripts denote time and superscripts denote arm. Thus, S_t^i is the state of arm i at time t . Bold letters denote the collection of variables across all arms. Thus, $\mathbf{S}_t = (S_t^1, \dots, S_t^n)$ is the set of states of all arms at time t . $S_{0:t}$ is a shorthand for (S_0, \dots, S_t) . We use $\mathbb{E}[\cdot]$ to denote the expectation of a random variable, $\mathbb{P}(\cdot)$ to denote probability of an event, and $\mathbb{1}\{\dots\}$ to denote the indicator of an event. Let $[n] := \{1, \dots, n\}$.

For a given function $f: \mathcal{X} \rightarrow \mathbb{R}$, the span norm of f is defined as $\text{span}(f) = \max_{x \in \mathcal{X}} f(x) - \min_{x \in \mathcal{X}} f(x)$. Given two metric spaces (\mathcal{X}, d_X) and (\mathcal{Y}, d_Y) , the Lipschitz constant of function $f: \mathcal{X} \rightarrow \mathcal{Y}$ is defined by

$$L_f = \sup_{\substack{x_1, x_2 \in \mathcal{X} \\ x_1 \neq x_2}} \frac{d_Y(f(x_1), f(x_2))}{d_X(x_1, x_2)}.$$

Let ζ_1 and ζ_2 denote probability measures on (\mathcal{X}, d_X) . Then, the Kantorovich distance between them is defined as

$$\mathcal{K}(\zeta_1, \zeta_2) = \sup_{f: L_f \leq 1} \left| \sum_{x \in \mathcal{X}} f(x) \zeta_1(x) - \sum_{x \in \mathcal{X}} f(x) \zeta_2(x) \right|.$$

Given a metric space (\mathcal{X}, d_X) , $\text{diam}(\mathcal{X}) = \sup\{d_X(x_1, x_2) : x_1, x_2 \in \mathcal{X}\}$ denotes the diameter of the set \mathcal{X} .

II. MODEL, PROBLEM FORMULATION, AND RESULTS

A. Restless Bandits

RBs are a class of resource allocation problems where, at each time instant, a decision maker has to select m out of n available alternatives. Each alternative, which is also called an arm, is a controlled Markov process $\langle \mathcal{S}^i, \mathcal{A}^i = \{0, 1\}, P^i, r^i \rangle$, where \mathcal{S}^i is the state space, $\mathcal{A}^i = \{0, 1\}$ is the action space, $P^i: \mathcal{S}^i \times \mathcal{A}^i \rightarrow \Delta(\mathcal{S}^i)$ is the controlled transition matrix, and $r^i: \mathcal{S}^i \times \mathcal{A}^i \rightarrow [0, R_{\max}]$ is the per-step reward. The action $A_t^i = 1$ means the decision maker selects arm i at time t . The arms for which $A_t^i = 1$ are called *active* arms and the arms for which $A_t^i = 0$ are called *passive* arms.

Let $\mathcal{S} = \mathcal{S}^1 \times \dots \times \mathcal{S}^n$ denote the joint state space and $\mathcal{A}(m) = \{\mathbf{a} \in \{0, 1\}^n : \|\mathbf{a}\|_1 = m\}$ denote the feasible action space, where $\|\mathbf{a}\|_1 := \sum_{i \in [n]} |a^i| = \sum_{i \in [n]} a^i$. We assume that the initial state $\mathbf{S}_0 = (S_0^1, \dots, S_0^n)$ is a random variable which is independent across arms and has a known initial distribution. Moreover, the arms evolve independently, i.e., for any $\mathbf{s}_{0:t} = (s_{0:t}^1, \dots, s_{0:t}^n)$ and $\mathbf{a}_{0:t} = (a_{0:t}^1, \dots, a_{0:t}^n)$,

we have

$$\begin{aligned} & \mathbb{P}(S_{t+1} = s_{t+1} | S_{0:t} = s_{0:t}, \mathbf{A}_{0:t} = \mathbf{a}_{0:t}) \\ &= \prod_{i=1}^n P^i(s_{t+1}^i | s_t^i, a_t^i) := P(s_{t+1} | s_t, \mathbf{a}_t). \end{aligned}$$

We consider two reward models as follows.

- 1) *Model A*: All arms, whether active or not, yield rewards, i.e., the aggregated per-step reward is given by $\mathbf{r}(s_t, \mathbf{a}_t) = \sum_{i \in [n]} r^i(s_t^i, a_t^i)$.
- 2) *Model B*: Only the activated arms yield rewards. The state of the passive arms evolves, but the arms do not yield reward. Thus, the aggregated per-step reward is given by $\mathbf{r}(s_t, \mathbf{a}_t) = \sum_{i \in [n]} r^i(s_t^i, a_t^i) \mathbb{1}(\{a_t^i = 1\})$.

Note that Model B is the same as Model A under the assumption that $r^i(\cdot, 0) = 0$ for all arms $i \in [n]$. For that reason, for most of the article, we will take $\mathbf{r}(s_t, \mathbf{a}_t) = \sum_{i \in [n]} r^i(s_t^i, a_t^i)$ and assume $r^i(\cdot, 0) = 0$ when specializing for results of Model B.

Remark 1: Both Models A and B arise in different applications. Examples of Model A include queuing networks [2], where all queues incur a holding cost, and machine maintenance [21], where all machines incur a cost when run in a faulty state. Examples of Model B include cognitive radios [8], where the reward depends only on the state of the selected channels.

Let Π denote the family of all possible (potentially history dependent and randomized) policies for the decision maker (who observes the state of all arms). The performance of any policy $\pi \in \Pi$ is given by

$$J(\pi) := \liminf_{T \rightarrow \infty} \frac{1}{T} \mathbb{E} \left[\sum_{t=1}^T \mathbf{r}(S_t, \mathbf{A}_t) \right] \quad (1)$$

where the expectation is taken with respect to the initial state distribution and the joint distribution induced on all system variables.

The objective of the decision maker is to choose a policy $\pi \in \Pi$ to maximize the total expected reward $J(\pi)$. This objective is an MDP but computing an optimal policy using a dynamic program suffers from the curse of dimensionality. For example, if $|\mathcal{S}^i| = S$ for each $i \in [n]$, then $|\mathcal{S}| = S^n$ and $|\mathcal{A}(m)| = \binom{n}{m}$. Then, the computational complexity of each step of value iteration is $|\mathcal{A}(m)| |\mathcal{S}|^2 = \binom{n}{m} S^{2n}$, which is prohibitively large for even moderate values of S and n . For this reason, most of the RB literature focuses on a computationally tractable but suboptimal approach known as the Whittle index policy.

B. Whittle Index Policy

The Whittle index policy is motivated by the solution of a relaxation of the original optimization problem. Instead of the hard constraint of activating exactly m arms at a time, consider a relaxation where m arms have to be activated on average, i.e.,

$$\begin{aligned} & \max_{\pi \in \Pi} \liminf_{T \rightarrow \infty} \frac{1}{T} \mathbb{E} \left[\sum_{t=1}^T \mathbf{r}(S_t, \mathbf{A}_t) \right], \\ & \text{s.t. } \limsup_{T \rightarrow \infty} \frac{1}{T} \mathbb{E} \left[\sum_{t=1}^T \|\mathbf{A}_t\|_1 \right] = m. \end{aligned} \quad (2)$$

Note that this relaxation is simply used to obtain a decomposition to define Whittle indices. The Whittle index policy, which is stated at the end of this section, picks exactly m arms at each time step.

Following [44], the Lagrangian relaxation of (2) is given by:

$$\mathcal{L}(\lambda, \pi) = \liminf_{T \rightarrow \infty} \frac{1}{T} \mathbb{E} \left[\sum_{t=1}^T \left[\mathbf{r}(S_t, \mathbf{A}_t) - \lambda \|\mathbf{A}_t\|_1 \right] \right]. \quad (3)$$

Moreover, a policy $\pi^* \in \Pi$ is optimal if there exists a Lagrange multiplier $\lambda \geq 0$ such that $\pi^* \in \arg \max_{\pi \in \Pi} \mathcal{L}(\lambda, \pi)$ and the expected number of activations under π^* is exactly m . See [44, Th. 4.3].

Note that the optimization problem $\max_{\pi \in \Pi} \mathcal{L}(\lambda, \pi)$ is decoupled across arms because the per-step reward is decoupled

$$\mathbf{r}(S_t, \mathbf{A}_t) - \lambda \|\mathbf{A}_t\|_1 = \sum_{i \in [n]} \left[r^i(S_t^i, A_t^i) - \lambda A_t^i \right].$$

Therefore, for a given λ , $\max_{\pi \in \Pi} \mathcal{L}(\lambda, \pi)$ is equivalent to the following n decoupled optimization problems: For all $i \in [n]$

$$\max_{\pi^i: \mathcal{S}^i \rightarrow \{0,1\}} \liminf_{T \rightarrow \infty} \frac{1}{T} \mathbb{E} \left[\sum_{t=1}^T \left[r^i(S_t^i, A_t^i) - \lambda A_t^i \right] \right]. \quad (4)$$

Let π_λ^i denote the optimal policy for Problem (4). Define the *passive set* \mathcal{W}_λ^i as the set of states for which the optimal policy π_λ^i prescribes passive action, i.e., $\mathcal{W}_\lambda^i := \{s \in \mathcal{S}^i : \pi_\lambda^i(s) = 0\}$.

Definition 1 (Indexability and Whittle index): An RB is said to be *indexable* if \mathcal{W}_λ^i is nondecreasing in λ , i.e., for any $\lambda_1, \lambda_2 \in \mathbb{R}$ such that $\lambda_1 \leq \lambda_2$, we have $\mathcal{W}_{\lambda_1}^i \subseteq \mathcal{W}_{\lambda_2}^i$. For an indexable RB, the *Whittle index* $w^i(s)$ of state $s \in \mathcal{S}^i$ is the smallest value of λ for which state s is part of the passive set \mathcal{W}_λ^i , i.e.,

$$w^i(s) = \inf \{ \lambda \in \mathbb{R} : s \in \mathcal{W}_\lambda^i \}.$$

Note that if the penalty $\lambda = w^i(s)$, then the policy π_λ^i is indifferent between taking passive or active actions at state s .

The Whittle index policy is a feasible policy for the original optimization problem and is given as follows: *At each time, activate the arms with the m largest values of the Whittle index at their current state.*

As argued in [16], the Whittle index policy is meaningful only when all arms are indexable. Various sufficient conditions for indexability are available in [21], [22], and [24]. In some settings, the Whittle index policy is optimal [6], [17], [18]. For general models, there is also strong evidence to suggest that the Whittle index policy performs close to optimal [21], [22], [24], [45], [46]. Algorithms to efficiently compute Whittle indices are presented in [24] and [47].

C. Learning Problem

Let μ denote the Whittle index policy and $J(\mu)$ denote its performance. We are interested in a setting where the transition matrices $\{P^i\}_{i \in [n]}$ of the arms are unknown but the decision maker has a prior on them. In this setting, the performance of a policy π operating for horizon T is characterized by the

Bayesian regret given by

$$\mathcal{R}(T; \pi) = \mathbb{E}^\pi \left[T\mathbf{J}(\boldsymbol{\mu}) - \sum_{t=1}^T \mathbf{r}(\mathbf{S}_t, \mathbf{A}_t) \right] \quad (5)$$

where the expectation is taken with respect to the prior distribution on $\{P^i\}_{i \in [n]}$, the initial condition, and the potential randomization done by the policy π . The Bayesian regret is a well-known metric used in various settings [43], [48], [49], [50]. An alternative method to quantify the performance is the frequentist regret but we focus on the Bayesian regret for a comparison of the two notions of the regret; see [51] and [52].

Remark 2: We measure the regret with respect to the Whittle index policy. In contrast, in most of the existing research, regret is defined with reference to the optimal policy. In principle, the results presented in this article are also applicable to the regret defined with respect to the optimal policy provided it is possible to compute the optimal policy for a given model. See Section VI-C for details.

Remark 3: The rested multiarmed bandit problem is a special case of Model B, where passive arms are frozen, i.e., $P^i(s_+ | s, 0) = \mathbb{1}\{s_+ = s\}$ for all arms $i \in [n]$. For this model, the Whittle index policy reduces to what is called the *Gittins index policy* and is optimal [17]. Thus, the results obtained in this article are also applicable to the rested multiarmed bandits.

D. Main Results

Our main result is to propose a Thompson-sampling-based algorithm, which we call RB-TSDE, and characterize its regret. In particular, let $S^i = |\mathcal{S}^i|$ denote the size of the state space of arm i and $\bar{S}_n = \sum_{i \in [n]} S^i$ denote the sum of the sizes of the state space of all arms. Then, we show the following.

Main Result: The regret of RB-TSDE is bounded by

$$\mathcal{R}(T; \text{RB-TSDE}) \leq \mathcal{O}(\alpha \bar{S}_n \sqrt{T \log T})$$

where $\alpha = n$ for Model A and $\alpha = m$ for Model B. Under additional assumptions, the bound for both models can be tightened to

$$\mathcal{R}(T; \text{RB-TSDE}) \leq \mathcal{O}(\max\{\alpha \sqrt{\bar{S}_n}, \bar{S}_n\} \sqrt{T \log T}).$$

The detailed characterization of the constants in the $\mathcal{O}(\cdot)$ terms is given in Theorems 1 and 2 later.

III. LEARNING ALGORITHM FOR RB

A. Assumptions on the Unknown Parameters

Let θ_*^i denote the unknown parameters of the transition matrices $[P^i(\cdot | \cdot, 0) P^i(\cdot | \cdot, 1)]$, $i \in [n]$. We assume that θ_*^i belongs to a compact set Θ^i . We impose the following assumptions on the model.

Assumption 1: For any $i \in [n]$ and $\theta^i \in \Theta^i$, the RB $\langle \mathcal{S}^i, \mathcal{A}^i = \{0, 1\}, P^i(\theta^i), r^i \rangle$ is indexable.

Assumption 2: Let $\boldsymbol{\mu}_\theta$ denote the Whittle index policy corresponding to model θ . Let \mathbf{P}_θ denote the controlled transition matrix under policy $\boldsymbol{\mu}_\theta$ and J_θ denote the average reward of policy $\boldsymbol{\mu}_\theta$. We assume that for every $\theta \in \Theta$, J_θ does not depend on the initial state and also assume that there exists a bounded

differential value function \mathbf{V}_θ such that $(J_\theta, \mathbf{V}_\theta)$ satisfy the average reward Bellman equation

$$J_\theta + \mathbf{V}_\theta(\mathbf{s}) = \mathbf{r}(\mathbf{s}, \boldsymbol{\mu}_\theta(\mathbf{s})) + [\mathbf{P}_\theta \mathbf{V}_\theta](\mathbf{s}), \quad \forall \mathbf{s} \in \mathcal{S}. \quad (6)$$

Assumption 1 is necessary for the Whittle index heuristic to be meaningful. Assumption 2 ensures that the average reward of the Whittle index policy is well defined for all models. There are various sufficient conditions for Assumption 2 in literature. See [53] for an overview.

The average reward Bellman equation (6) has an infinite number of solutions. In particular, if $(J_\theta, \mathbf{V}_\theta)$ satisfies (6), then so does $(J_\theta, \mathbf{V}_\theta + \text{constant})$. Assumption 2 implies that $\text{span}(\mathbf{V}_\theta)$ is bounded. A bound on $\text{span}(\mathbf{V}_\theta)$ under a different set of assumptions is presented in [39] but this bound does not suffice for our analysis. See Remark 7 for details. As we want to capture the scaling of $\text{span}(\mathbf{V}_\theta)$ with n and m , so we impose an additional assumption on the model.

The ergodicity coefficient of \mathbf{P}_θ is defined as

$$\lambda_{\mathbf{P}_\theta} = 1 - \min_{\substack{\mathbf{s}, \mathbf{s}' \in \mathcal{S} \\ \mathbf{z} \in \mathcal{S}}} \min\{\mathbf{P}_\theta(\mathbf{z} | \mathbf{s}), \mathbf{P}_\theta(\mathbf{z} | \mathbf{s}')\}.$$

We impose the following assumption on the model.

Assumption 3: We assume there exists $\lambda^* < 1$ such that $\sup_{\theta \in \Theta} \lambda_{\mathbf{P}_\theta} \leq \lambda^*$.

See [53, Sec. 5] for various equivalent characterizations of $\lambda_{\mathbf{P}_\theta} < 1$. Assumption 3 is used while analyzing the rate of convergence of relative value iteration [54] to bound the span of the value function. A sufficient condition of Assumption 3 is presented in [55, Th. 8.11] and a relaxation of the assumption is presented in Section VI-A.

The ergodicity coefficient of \mathbf{P}_θ depends on the Whittle index policy $\boldsymbol{\mu}_\theta$. A policy independent upper bound of the ergodicity coefficient is given by the contraction factor, which is defined as

$$\lambda' = 1 - \min_{\substack{\mathbf{s}, \mathbf{s}' \in \mathcal{S}, \\ \mathbf{a} \in \mathcal{A}(m), \mathbf{a}' \in \mathcal{A}(m)}} \sum_{\mathbf{z} \in \mathcal{S}} \min\{\mathbf{P}(\mathbf{z} | \mathbf{s}, \mathbf{a}), \mathbf{P}(\mathbf{z} | \mathbf{s}', \mathbf{a}')\}.$$

Since the dynamics of the arms are independent, the definition of contraction factor implies that a sufficient condition for Assumption 3 is that for every arm, and every pair of state-action pairs, there exists a next state that can be reached from both state-action pairs with positive probability in one step. Moreover, if for every arm there is a distinguished state which can be reached from all state-action pairs with probability at least ε , then the ergodicity coefficient is less than $1 - \varepsilon$.

B. Priors and Posterior Updates

We assume that $\{\theta_*^i\}_{i \in [n]}$ are independent random variables and use ϕ_1^i to denote the prior on θ^i for each arm $i \in [n]$. At time t , let $h_t^i = (s_1^i, a_1^i, \dots, s_{t-1}^i, a_{t-1}^i, s_t^i)$ denote the history of states and actions at arm i . Let ϕ_t^i denote the posterior distribution on θ_*^i given h_t^i . Then, upon applying action a_t^i and observing the next state s_{t+1}^i , the posterior distribution ϕ_{t+1}^i can be computed using the Bayes rule as

$$\phi_{t+1}^i(d\theta) = \frac{P^i(s_{t+1}^i | s_t^i, a_t^i; \theta) \phi_t^i(d\theta)}{\int P^i(s_{t+1}^i | s_t^i, a_t^i; \tilde{\theta}) \phi_t^i(d\tilde{\theta})}. \quad (7)$$

Algorithm 1: RB-TSDE.

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1: Input:  $\{(\Theta^i, \phi_1^i)\}_{i \in [n]}$ .
2: Initialize  $t \leftarrow 1, t_1 \leftarrow 0, T_0 \leftarrow 0, k \leftarrow 0,$ 
    $N^i(s^i, a^i) = 0, \forall a^i \in \{0, 1\}, s^i \in \mathcal{S}^i, \forall i \in [n], \theta_0, \mu_{\theta_0}.$ 
3: for  $t = 1, 2, \dots$  do
4:   if  $t_k + T_{k-1} < t$  or  $2N_{t_k}^i(s, a) < N_t^i(s, a)$  for any
      $s \in \mathcal{S}^i, a \in \{0, 1\}, i \in [n]$  then
5:     Set  $k \leftarrow k + 1, T_{k-1} \leftarrow t - t_{k-1}, t_k \leftarrow t.$ 
6:     Actor  $i \in [n]$  samples  $\theta_{t_k}^i \sim \phi_{t_k}^i$  and compute  $w_{t_k}^i.$ 
7:   end if
8:   Actor  $i \in [n]$  sends the Whittle index  $w_{t_k}^i(s_t^i)$  to the
     coordinator.
9:   The coordinator sends  $a_t^i = 1$  to the arms with the
      $m$ -highest Whittle index and sends  $a_t^i = 0$  to others.
10:  Actor  $i \in [n]$  updates  $\phi_{t+1}^i$  according to (7).
11: end for

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If the prior is a conjugate distribution on Θ^i , then the posterior can be updated in closed form. Note that the exact form of the prior and the posterior update rule are not important for the description of the algorithm or its regret analysis.

C. RB-TSDE Algorithm (Distributed Implementation)

We propose a Thompson-sampling-based algorithm, which we call RB-TSDE. Our algorithm is inspired by the TSDE algorithm in [43].

The RB-TSDE algorithm consists of a coordinator and n actors, one for each arm. The coordinators and the actors require synchronized communication as described below. The whole algorithm is described in Algorithm 1.

As the name suggests, RB-TSDE operates in episodes of dynamic length. The episodes are synchronized for all actors and the coordinator signals the start of episodes to all actors. The actor at arm i maintains a posterior ϕ_t^i distribution on the dynamics of arm i according to (7) and keep track of $N_t^i(s^i, a^i) = \sum_{\tau=1}^{t-1} \mathbb{1}\{(S_\tau^i, A_\tau^i) = (s^i, a^i)\}.$

Let t_k and T_k denote the start time and length of episode k , respectively. The end of the episode can either be triggered by the coordinator or any of the actors. The coordinator triggers the end of the episode if the length of the episode is one more than the length of the previous episode. The actor for arm i triggers the end of the episode if the number of state-action visits $N_{t_k}^i(s_t^i, a_t^i)$ of the current state-action pair are more than double of their value at the beginning of the episode. Thus

$$t_{k+1} = \min\{t > t_k : t - t_k > T_{k-1} \text{ or } N_{t_k}^i(s^i, a^i) > 2N_{t_k}^i(s^i, a^i) \text{ for some } (i, s^i, a^i)\}.$$

At the beginning of episode k , the actor for arm $i \in [n]$ samples a parameter θ_k^i from the posterior $\phi_{t_k}^i$ and computes the Whittle index $w_{t_k}^i$ for all states. During episode k , at each time t , the actor at arm i sends the value of $w_{t_k}^i(s_t^i)$ to the coordinator. The coordinator receives $w_{t_k}^i(s_t^i)$ from all arms, sends the active action $a_t^i = 1$ to the arms with the m -highest values of the Whittle index, and sends the passive action $a_t^i = 0$ to the remaining

arms. This process continues until a condition for ending the episode is triggered by the coordinator or one of the actors.

D. Regret Bound

Theorem 1: Under Assumptions 1–3, the regret of RB-TSDE is upper bounded as follows:

$$\mathcal{R}(T; \text{RB-TSDE}) < 40\alpha \frac{R_{\max}}{1 - \lambda^*} \bar{S}_n \sqrt{T \log T}$$

where $\alpha = n$ for Model A and $\alpha = m$ for Model B.

See Section IV-D for proof.

We derive a tighter regret bound under a stronger assumption. We first assume that the state space of each arm \mathcal{S}^i , $i \in [n]$, is a finite subset of \mathbb{R} and use d^i to denote the Euclidean metric on \mathbb{R} , i.e., $d^i(s, s') = |s - s'|$. Furthermore, let $\mathbf{d}(s, s') = \sum_{i \in [n]} d^i(s^i, s'^i)$ for any $s, s' \in \mathcal{S}$. We then impose the following assumption.

Assumption 4: For each $\theta \in \Theta$, the value function V_θ is Lipschitz with a Lipschitz constant upper bounded by L_v .

In general, Assumption 4 depends on the specific model being considered. We present one instance where Assumption 4 is satisfied in Section VI-B.

Theorem 2: Under Assumptions 1–4, the regret of RB-TSDE for both Model A and Model B is upper bounded as follows:

$$\begin{aligned} \mathcal{R}(T; \text{RB-TSDE}) \\ < 16 \max\left\{ \alpha R_{\max} \sqrt{\bar{S}_n}, D_{\max} L_v \bar{S}_n \right\} \\ \sqrt{KT \log(T \max\{1, K'\})} \end{aligned}$$

where $\alpha = n$ for Model A and $\alpha = m$ for Model B, K , and K' are positive constants independent of n and T , and $D_{\max} = \max_{i \in [n]} \text{diam}(\mathcal{S}^i)$.

See Section IV-D for proof.

Remark 4: If we directly use an existing RL algorithm for RBs, the regret will scale as $\tilde{O}(2^n \sqrt{T})$ or larger. The results of Theorems 1 and 2 show that the regret scales as either $\tilde{O}(n^2 \sqrt{T})$, $\tilde{O}(n^{1.5} \sqrt{T})$, or $\tilde{O}(n \sqrt{T})$ depending on the modeling assumptions. Thus, using a learning algorithm which is adapted to the structure of the models gives a significantly better scaling with the number of arms.

Remark 5: The exact scaling with the number of arms depends on how m scales with n . For example, if m remains constant, then under Assumptions 1–3, the regret for Model A scales as $\tilde{O}(n^2 \sqrt{T})$, while the regret for Model B scales as $\tilde{O}(n \sqrt{T})$. Assumptions 1–3, if L_v is upper bounded by a constant, under Assumption 4, the regret for Model A scales as $\tilde{O}(n^{1.5} \sqrt{T})$, while the regret for Model B scales as $\tilde{O}(n \sqrt{T})$. On the other hand, if m scales as βn , where $\beta < 1$ is a constant,¹ then under Assumptions 1–3, the regret for both models scales as $\tilde{O}(n^2 \sqrt{T})$. Under Assumption 4, the regret for both models scales as $\tilde{O}(n^{1.5} \sqrt{T})$. Thus, for Model A, the regret bound of Theorem 2 is tighter than that of Theorem 1, but for Model B, it depends on the scaling assumptions on m .

¹For this setting, it was shown in [18] that the Whittle index policy is optimal as $n \rightarrow \infty$.

IV. REGRET ANALYSIS

The high-level idea of the proof is similar to the analysis in [43], but we exploit the properties of the RB model while simplifying individual terms. We first start with bounds on the average reward and the differential value function.

A. Bounds on Average Reward and Differential Value Function

As mentioned earlier, \mathbf{V}_θ is unique only up to an additive constant. We assume that \mathbf{V}_θ is chosen such that $\xi_\theta \mathbf{V}_\theta = 0$, where ξ_θ is the stationary distribution of \mathbf{P}_θ . This ensures that \mathbf{V}_θ is equal to the *asymptotic bias* of policy μ_θ and is given by

$$\mathbf{V}_\theta = \sum_{t=0}^{\infty} \mathbf{P}_\theta^t [\mathbf{r} - J_\theta]. \quad (8)$$

See for example [56].

Then we have the following bounds.

Lemma 1: Under Assumption 3, for any $\theta \in \Theta$,

$$0 \leq J_\theta \leq \alpha R_{\max} \text{ and } \text{span}(\mathbf{V}_\theta) \leq 2\alpha R_{\max}/(1 - \lambda^*)$$

where $\alpha = n$ for Model A and $\alpha = m$ for Model B.

Proof: Note that for Model A, $\mathbf{r}(s, \mathbf{a}) \in [0, nR_{\max}]$ while for Model B, $\mathbf{r}(s, \mathbf{a}) \in [0, mR_{\max}]$. Then, the bounds for J_θ follow immediately from definition. The bounds on $\text{span}(\mathbf{V}_\theta)$ follow from (8), $\text{span}(s + \mathbf{y}) \leq \text{span}(s) + \text{span}(\mathbf{y})$, Assumption 3, and the fact that for any vector \mathbf{v} , $\text{span}(\mathbf{P}_\theta \mathbf{v}) \leq \lambda_{\mathbf{P}_\theta} \text{span}(\mathbf{v})$. ■

Remark 6: Lemma 1 shows the key difference between Models A and B. When all arms yield rewards, the maximum value of $\mathbf{r}(s, \mathbf{a})$ is nR_{\max} while when only active arms yield rewards, the maximum value of $\mathbf{r}(s, \mathbf{a})$ is mR_{\max} . This leads to different bounds on J_θ and \mathbf{V}_θ .

Remark 7: An alternative bound on $\text{span}(\mathbf{V}_\theta)$ is presented in [39]. Let $T_\theta^{s_1 \rightarrow s_2}$ denote the expected number of steps to go from state s_1 to state s_2 under policy μ_θ for model θ . Define $\mathbf{D}_\theta = \max_{s_1, s_2} T_\theta^{s_1 \rightarrow s_2}$ to be the *one-way diameter*. Then, it is shown in [39] that $\text{span}(\mathbf{V}_\theta) \leq J_\theta \mathbf{D}_\theta$. We do not know of an easy way to characterize the dependence of \mathbf{D}_θ on the number n of arms. That is why we consider an alternative bound on $\text{span}(\mathbf{V}_\theta)$.

B. Regret Decomposition

For the ease of notation, we simply use $\mathcal{R}(T)$ instead of $\mathcal{R}(T; \text{RB-TSDE})$. We also use $(J_*, \mu_*, \mathbf{P}_*, \mathbf{V}_*)$ instead of $(J_{\theta_*}, \mu_{\theta_*}, \mathbf{P}_{\theta_*}, \mathbf{V}_{\theta_*})$ and use $(J_k, \mu_k, \mathbf{P}_k, \mathbf{V}_k)$ instead of $(J_{\theta_k}, \mu_{\theta_k}, \mathbf{P}_{\theta_k}, \mathbf{V}_{\theta_k})$. Rearranging terms in Bellman (6) and adding and subtracting $\mathbf{V}_k(s_{t+1})$, we get

$$\begin{aligned} \mathbf{r}(s_t, \mathbf{a}_t) &= J_k + \mathbf{V}_k(s_t) - \mathbf{V}_k(s_{t+1}) + \mathbf{V}_k(s_{t+1}) \\ &\quad - [\mathbf{P}_k \mathbf{V}_k](s_t). \end{aligned} \quad (9)$$

Let K_T denote the number of episodes until horizon T . Substituting (9) in (5), we get

$$\begin{aligned} \mathcal{R}(T) &= \underbrace{\mathbb{E} \left[T J_* - \sum_{k=1}^{K_T} T_k J_k \right]}_{\text{regret due to sampling error} =: \mathcal{R}_0(T)} \\ &\quad + \underbrace{\mathbb{E} \left[\sum_{k=1}^{K_T} \sum_{t=t_k}^{t_{k+1}-1} \mathbf{V}_k(s_{t+1}) - \mathbf{V}_k(s_t) \right]}_{\text{regret due to time-varying policy} =: \mathcal{R}_1(T)} \\ &\quad + \underbrace{\mathbb{E} \left[\sum_{k=1}^{K_T} \sum_{t=t_k}^{t_{k+1}-1} [\mathbf{P}_k \mathbf{V}_k](s_t) - \mathbf{V}_k(s_{t+1}) \right]}_{\text{regret due to model mismatch} =: \mathcal{R}_2(T)}. \end{aligned} \quad (10)$$

C. Bounding Individual Terms

Each term of (10) is bounded as follows.

Lemma 2: Under Assumptions 1–3, we have

- 1) $\mathcal{R}_0(T) \leq 2\alpha R_{\max} \sqrt{\bar{S}_n T \log T}$;
- 2) $\mathcal{R}_1(T) \leq 4 \frac{\alpha R_{\max}}{1 - \lambda^*} \sqrt{\bar{S}_n T \log T}$;
- 3) $\mathcal{R}_2(T) \leq 12\sqrt{2} \frac{\alpha R_{\max}}{1 - \lambda^*} (n + \bar{S}_n \sqrt{T \log T})$.

See the Appendix for the proof steps.

We can obtain an alternative bound on $\mathcal{R}_2(T)$ under Assumption 4.

Lemma 3: Under Assumptions 1, 2, and 4, we have

- 1) $\mathcal{R}_0(T) \leq 2\alpha R_{\max} \sqrt{\bar{S}_n T \log T}$.
- 2) $\mathcal{R}_1(T) \leq 2D_{\max} \mathbf{L}_v \sqrt{\bar{S}_n T \log T}$.
- 3) $\mathcal{R}_2(T) \leq 12\bar{S}_n D_{\max} \mathbf{L}_v \sqrt{KT \log(K'T)}$

where K and K' are positive constants that do not depend on n and T .

The bound on $\mathcal{R}_0(T)$ is similar as Lemma 3. The bound on $\mathcal{R}_1(T)$ is established based on Appendix A and the fact that Lipschitz continuity of \mathbf{V}_θ implies that $\text{span}(\mathbf{V}_\theta) \leq \mathbf{L}_v \text{diam}(\mathcal{S})$. The bound on $\mathcal{R}_2(T)$ is obtained by following the steps in Appendix D.

D. Obtaining the Final Bound

Proof of Theorem 1: From (10) and Lemma 2, we get

$$\begin{aligned} \mathcal{R}(T) &\leq 2\alpha R_{\max} \sqrt{\bar{S}_n T \log T} + 4 \frac{\alpha R_{\max}}{1 - \lambda^*} \sqrt{\bar{S}_n T \log T} \\ &\quad + 12\sqrt{2} \frac{\alpha R_{\max}}{1 - \lambda^*} (n + \bar{S}_n \sqrt{T \log T}). \end{aligned}$$

By definition, $\lambda^* < 1$. Then

$$\begin{aligned} \mathcal{R}(T) &< 6 \frac{\alpha R_{\max}}{1 - \lambda^*} \sqrt{\bar{S}_n T \log T} + 24\sqrt{2} \frac{\alpha R_{\max}}{1 - \lambda^*} \bar{S}_n \sqrt{T \log T} \\ &< (6 + 24\sqrt{2}) \frac{\alpha R_{\max}}{1 - \lambda^*} \bar{S}_n \sqrt{T \log T} \\ &= 40 \frac{\alpha R_{\max}}{1 - \lambda^*} \bar{S}_n \sqrt{T \log T}. \end{aligned}$$

This completes the proof of Theorem 1. ■

Proof of Theorem 2: From (10), Lemma 2, and 3, we get

$$\begin{aligned} \mathcal{R}(T) &\leq 2\alpha R_{\max} \sqrt{\bar{S}_n T \log T} + 2D_{\max} L_v \sqrt{\bar{S}_n T \log T} \\ &\quad + 12\bar{S}_n D_{\max} L_v \sqrt{KT \log(K'T)} \\ &< 2\alpha R_{\max} \sqrt{\bar{S}_n T \log T} \\ &\quad + 14\bar{S}_n D_{\max} L_v \sqrt{KT \log(K'T)}. \end{aligned}$$

Let $\bar{R}_n = \max\{\alpha R_{\max}/\sqrt{\bar{S}_n}, D_{\max} L_v\}$. Then, we have

$$\mathcal{R}(T) < 16\bar{S}_n \bar{R}_n \sqrt{KT \log(K'T)}.$$

This completes the proof of Theorem 2. \blacksquare

V. NUMERICAL EXAMPLES

In this section, we demonstrate the empirical performance of RB-TSDE. In particular, we consider two environments, one for Model A (a model for machine maintenance) and one for Model B (a model for link scheduling). For both environments, we consider multiple experiments and plot the regret as a function of time and as a function of number of arms. Our results illustrate that the regret does indeed scale according to our theoretical results. We also compare the results with the empirical performance of QWI, which is a Q-learning algorithm for RBs proposed in [26], [28], and [29]. Note that only the algorithm proposed in [29] is called QWI, but the algorithms in [26] and [28] are conceptually similar, so we collectively call them QWI.

A. Environments

We start with a description of the two environments.

1) Environment A: We consider a machine maintenance model where a single repairman is responsible for the maintenance of a set of machines, which deteriorate over time. Each machine has multiple deterioration states sorted from *pristine* to *ruined*. There is a cost associated with running the machine and the cost is nondecreasing function of the state. If the machine is left unmonitored, then the state of the machine stochastically deteriorates over time. The repairman may visit one of the machines and replace it with a new machine at a fixed cost. The objective is to determine a scheduling policy to minimize the expected discounted cost over time.

We model the above environment as an instance of Model A. In particular, we consider n arms, where $n \in \{10, 20, \dots, 80\}$ where $m = 1$ arm can be activated at each time. The state space of each arm is of size $S = 10$. Under $a = 1$, the state of the arm is reset to 1 (and this fact is known to the learner). Under $a = 0$, the transition matrix is stochastic monotone and chosen, as described in [24, Appendix 1.2]. The transitions under the passive action are unknown to the learner. The per-step reward function are given by $r^i(s, 0) = (S - 1)^2 - (s - 1)^2$, $r^i(s, 1) = 0.5(S - 1)^2$ for all $i \in [n]$ and $s \in [S]$.

2) Environment B: We consider a link scheduling problem where there are n users who can communicate over a shared communication link. Each user has a queue, where packets arrive

according to an unknown i.i.d. process. At each time, a controller may schedule one of the users and transmit all its packets over the channel. The users which are not scheduled incur a holding cost which is equal to the square of the number of packets in the queue.

We model the above environment as an instance of Model B. In particular, we consider n arms, where $n \in \{10, 20, \dots, 80\}$ where $m = n - 1$ arm can be activated at each time. The state of each arm is of size $S = 10$. Under $a = 1$, the transition matrix is upper triangular and chosen, as described in [57, $\mathcal{P}_1(p)$ of Appendix A], where p is set to be different for each arm, linearly ranged from 0.05 to 0.95. The transitions under the active action are unknown to the learner. Under $a = 0$, the state of the arm is reset to 1 (and this fact is known to the learner). The per-step reward function are given by $r^i(s, 0) = 0$, $r^i(s, 1) = (S - 1)^2 - (s - 1)^2$ for all $i \in [n]$ and $s \in [S]$.

B. Algorithms

We compare the performance of two algorithms.

1) RB-TSDE: We consider the RB-TSDE algorithm described in Algorithm 1. We initialize the algorithm with un-informed Dirichlet prior on the unknown parameters and update the posterior according to the conjugate posterior for Dirichlet priors.

2) QWI: We also consider QWI, which is a Q-learning algorithm for RBs proposed in [26], [28], and [29] as a baseline. The algorithm has two learning rates. As recommended in [28, eq. (17)] we pick the step-size sequence which has a good performance by setting parameters C and C' of QWI as $C = 0.03$ and $C' = 0.01$, where the numerical values were obtained by running a hyperparameter search.

C. Experimental Results

In our experiments, we pick a horizon of $T = 5000$ and compute the Bayesian regret averaged over 250 sample paths. We repeat the experiment for $n \in \{10, 20, \dots, 80\}$. For each environment, we plot four curves: (a) plot of $\mathcal{R}(T)/\sqrt{T}$ vs T for RB-TSDE; (b) plot of $\mathcal{R}(T)$ at $T = 5000$ versus n for RB-TSDE; (c) plot of $\mathcal{R}(T)/\sqrt{T}$ vs T for QWI; and (d) plot of $\mathcal{R}(T)$ at $T = 5000$ versus n for QWI. For plots (b) and (d), we also fit the points with a parametric curve of the form $p_0 + p_1 n + p_2 n^{1.5}$ for Environment A and $p_0 + p_1 n$ for Environment B to obtain the scaling with number of arms.

The plots for Environment A is shown in Fig. 1. The subplot (a) shows that the regret essentially scales as \sqrt{T} with time. The subplots (b) show that the regret scales as $n^{1.5}$ with the number of arms. Thus, the results are consistent with regret bounds of Theorem 2.

The plots for Environment B is shown in Fig. 2. The behavior of subplots (a) and (c) is the same as for Environment A. Note that for larger values of T , the plot of $\mathcal{R}(T)/\sqrt{T}$ has not yet converged to a straight line, but it is clear that these curves are upper bounded by a constant. The subplot (b) shows that the regret scales linearly with the number of arms. Note that in this case since $m = n - 1$, Theorem 1 suggests that the regret should

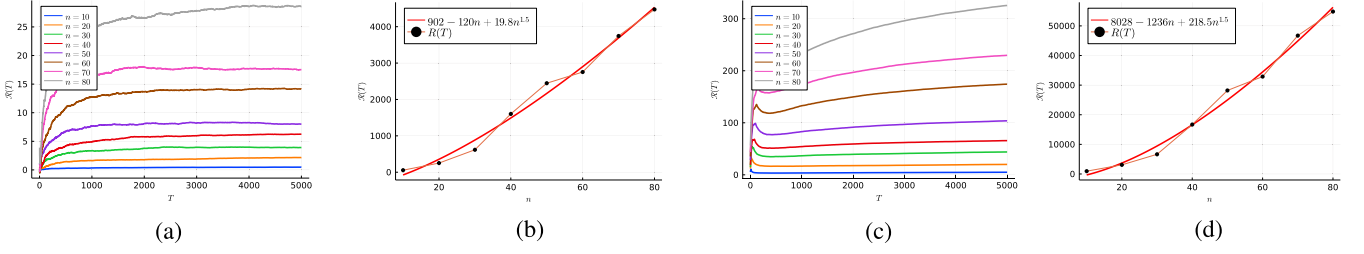


Fig. 1. Regret analysis of RB-TSDE and QWI method for Environment B. Note that RB-TSDE has an order of magnitude lower regret than QWI. (a) RB-TSDE: $\mathcal{R}(T)/\sqrt{T}$ versus T . (b) RB-TSDE: $\mathcal{R}(T)$ versus n . (c) QWI: $\mathcal{R}(T)/\sqrt{T}$ versus T . (d) QWI: $\mathcal{R}(T)$ versus n .

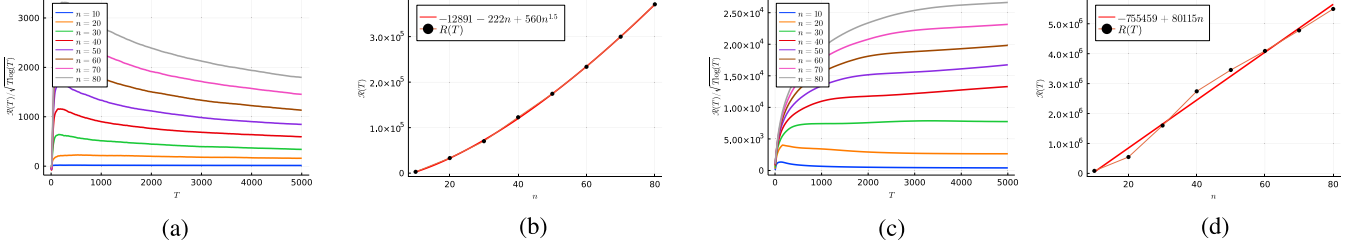


Fig. 2. Regret analysis of RB-TSDE and QWI for Environment A. Note that RB-TSDE has an order of magnitude better regret than QWI. (a) RB-TSDE: $\mathcal{R}(T)/\sqrt{T}$ versus T . (b) RB-TSDE: $\mathcal{R}(T)$ versus n . (c) QWI: $\mathcal{R}(T)/\sqrt{T}$ versus T . (d) QWI: $\mathcal{R}(T)$ versus n .

scale as $n^{1.5}$ (see Remark 5) and the result is consistent with the theorem.

Note that even though no regret analysis of the QWI algorithm for RBs is presented in [26], [28], and [29], the above experiment suggests that empirically the performance of QWI has similar features as RB-TSDE. However, unlike the QWI, RB-TSDE does not require any hyperparameter tuning. Moreover, RB-TSDE has an order of magnitude lower regret than QWI.

VI. DISCUSSION

A. Relaxation of Assumption 3

Assumption 3 can be relaxed as follows.

Assumption 5: For every θ , there exists a positive integer τ^* and a real $\lambda^* \in (0, 1)$ such that $\lambda_{P_{\theta}^*} \leq \lambda^*$.

Based on this assumption, the result of Lemmas 1 changes as follows.

Lemma 4: Under Assumption 5, for any $\theta \in \Theta$, $\text{span}(\mathbf{V}_{\theta}) \leq 2\tau^*\alpha R_{\max}/(1 - \lambda^*)$.

Consequently, we have the following changes.

Lemma 5: Under Assumptions 1, 2, and 5, we have

- 1) $\mathcal{R}_1(T) \leq 4 \frac{\tau^*\alpha R_{\max}}{1 - \lambda^*} \sqrt{\bar{S}_n T \log T}$;
- 2) $\mathcal{R}_2(T) \leq 12\sqrt{2} \frac{\tau^*\alpha R_{\max}}{1 - \lambda^*} (n + \bar{S}_n \sqrt{T \log T})$.

Theorem 3: Under Assumptions 1, 2, and 5, the regret of RB-TSDE is upper bounded as follows:

$$\mathcal{R}(T; \text{RB-TSDE}) < 40\alpha \frac{\tau^* R_{\max}}{1 - \lambda^*} \bar{S}_n \sqrt{T \log T}.$$

The proof steps of Lemmas 4 and 5 are similar to the proof steps of Lemmas 1 and 2; and the proof steps of Theorem 3 are similar to the proof steps of Theorems 1. See [58] for the details.

B. Set of Sufficient Conditions for Assumption 4

Assumption 6: Suppose each arm $i \in [n]$ is (L_r^i, L_p^i) Lipschitz, i.e.,

$$L_r^i = \sup_{s_{(1)}^i, s_{(2)}^i, a} \frac{|r^i(s_{(1)}^i, a) - r^i(s_{(2)}^i, a)|}{d^i(s_{(1)}^i, s_{(2)}^i)}$$

$$L_p^i = \sup_{s_{(1)}^i, s_{(2)}^i, a} \frac{\mathcal{K}(P^i(\cdot|s_{(1)}^i, a), P^i(\cdot|s_{(2)}^i, a))}{d^i(s_{(1)}^i, s_{(2)}^i)}$$

where $L_r^i < \infty$ and $L_p^i < 1$.

Assumption 7: For all $\theta \in \Theta$, the Whittle index policy is optimal.

Assumption 7 is satisfied in some instances, such as: i) the rested multiarmed bandit setup described in Remark 3, where only one arm can be activated at a time (i.e., $m = 1$) and yield reward, and arms that are not activated remain frozen (i.e., $P(s_+^i | s^i, 0) = \delta_{s^i}(s_+^i)$, where δ_{s^i} is the Dirac delta measure centered at s^i) [17]; ii) the number of arms are asymptotically large [18]; iii) certain queuing models [6].

Moreover, we assume that the product measure on \mathcal{S} is $d(s, s') = \sum_{i \in [n]} d^i(s^i, s'^i)$.

Lemma 6 ([59, Lemma 2]): Under Assumption 6, the MDP $\langle \mathcal{S}, \mathcal{A}(m), P, \mathcal{R} \rangle$ is $(\max_{i \in [n]} L_r^i, \max_{i \in [n]} L_p^i)$ -Lipschitz, i.e.,

$$\max_{\substack{s, s' \in \mathcal{S} \\ a \in \mathcal{A}(m)}} \frac{|r(s, a) - r(s', a)|}{d(s, s')} \leq \max_{i \in [n]} L_r^i$$

$$\max_{\substack{s, s' \in \mathcal{S} \\ a \in \mathcal{A}(m)}} \frac{\mathcal{K}(P(\cdot|s, a) - P(\cdot|s', a))}{d(s, s')} \leq \max_{i \in [n]} L_p^i.$$

An immediate consequence of Lemma 6 is the following.

Lemma 7: Under Assumptions 6 and 7, V_θ is Lipschitz with the Lipschitz constant bounded by

$$L_v \leq (\max_{i \in [n]} L_r^i) / (1 - \max_{i \in [n]} L_p^i).$$

Thus, Assumptions 6 and 7 imply Assumption 4.

Proof: The result follows from Lemma 6 and [60, Th. 4.2].

C. Regret With Respect to the Optimal Policy

We measure regret with respect to the Whittle index policy. For models where Assumption 7 is satisfied, the Whittle index policy is an optimal policy. Even when the assumption is not satisfied, it is possible to generalize the results of this article to identify the regret with respect to the optimal policy. In particular, let ψ_θ^* denote the optimal policy for model $\theta \in \Theta$. Then, the Bayesian regret of a learning algorithm π with respect to the optimal policy is

$$\mathcal{R}^*(T; \pi) = \mathbb{E}^\pi \left[T \mathbf{J}(\psi_\theta^*) - \sum_{t=1}^T \mathbf{r}(\mathbf{S}_t, \mathbf{A}_t) \right]. \quad (11)$$

Then, in principle, we can replace the distributed implementation presented in Algorithm 1 with a modified centralized implementation where the learner observes the state of all arms and maintains the posterior ϕ_t^i for all $i \in [n]$. At the beginning of each episode, the learner samples $\theta_{t_k}^i$ from $\phi_{t_k}^i$, computes the policy μ_{t_k} , which is optimal for the sampled model $(\theta_{t_k}^1, \dots, \theta_{t_k}^n)$, and plays μ_{t_k} for the rest of the episode. The regret of this variant will be the same as the bounds in Theorems 1 and 2. However, we do not present such an analysis here because it makes the resulting algorithm impractical as computing the optimal policy is intractable when there are more than a few arms.

VII. CONCLUSION

In this article, we present a Thompson-sampling-based RL algorithm for RBs. We show that the Bayesian regret of our algorithm with respect to an oracle that applies the Whittle index policy of the true model is either $\tilde{O}(nm\sqrt{T})$, $\tilde{O}(n^2\sqrt{T})$, $\tilde{O}(n^{1.5}\sqrt{T})$ or $\tilde{O}(n\sqrt{T})$ depending on assumptions on the model. These are in contrast to naively using any standard RL algorithm, which will have a regret that scales exponentially in n . Our results are also applicable to the rested multiarmed bandit setting, where the Whittle index policy is the same as the Gittins index and is optimal. All in all, our results illustrate that a learning algorithm which leverages the structure of the model can significantly improve regret compared to model-agnostic algorithms.

APPENDIX A

A. Bound on $\mathcal{R}_0(T)$ (Lemma 2.1)

We first state a basic property of the Thompson sampling algorithms.

Lemma 8 (TS Lemma [50]): Suppose the true parameters θ and the estimated ones θ_k have the same distribution given the

same history \mathcal{H} . For any \mathcal{H} -measurable function f , we have

$$\mathbb{E}[f(\theta)|\mathcal{H}] = \mathbb{E}[f(\theta_k)|\mathcal{H}].$$

Proof of Lemma 2.1: Now we consider $\mathcal{R}_0(T)$. Let $\tilde{J}_\star = \alpha R_{\max} - J_\star$ and $\tilde{J}_k = \alpha R_{\max} - J_k$. By Lemma 1, we have that $\tilde{J}_\star, \tilde{J}_k \in [0, R_{\max}]$. Therefore,

$$\begin{aligned} \mathcal{R}_0(T) &= \mathbb{E} \left[T J_\star - \sum_{k=1}^{K_T} T_k J_k \right] = \mathbb{E} \left[\sum_{k=1}^{K_T} T_k \tilde{J}_k - T \tilde{J}_\star \right] \\ &\stackrel{(a)}{\leq} \sum_{k=1}^{\infty} \mathbb{E} \left[\mathbb{1}(\{t_k \leq T\}) (T_{k-1} + 1) \tilde{J}_\star \right] - T \mathbb{E}[\tilde{J}_\star] \\ &\leq \mathbb{E} \left[\sum_{k=1}^{K_T} (T_{k-1} + 1) \tilde{J}_\star \right] - T \mathbb{E}[\tilde{J}_\star] \stackrel{(b)}{\leq} \alpha R_{\max} \mathbb{E}[K_T] \end{aligned}$$

where (a) uses the TS Lemma and the fact that due to the first stopping criterion, $T_k \leq T_{k-1} + 1$; (b) uses Lemma 1 and the fact that $\sum_{k=1}^{K_T} T_{k-1} \leq T$. The results follows from Lemma 9.

Lemma 9: The number of episodes K_T is bounded as follows:

$$K_T \leq 2\sqrt{\bar{S}_n T \log(T)}.$$

Proof: Define macro episodes with start times t_{n_l} , $l = 1, 2, \dots$ with $t_{n_1} = t_1$ and

$$\begin{aligned} t_{n_{l+1}} &= \min\{t_k > t_{n_l} : N_{t_k}^i(s^i, a^i) > 2N_{t_{k-1}}^i(s^i, a^i) \\ &\text{for some } (i, s^i, a^i)\}. \end{aligned} \quad (12)$$

Let γ be the number of macro episodes until time T and define $n_{(\gamma+1)} = K_T + 1$. The rest of the proof is the same as [43, eq. (8) in proof of Lemma 1] by which we get $K_T \leq \sqrt{2\gamma T}$.

For each arm-state-action tuple, define

$$\gamma^i(s^i, a^i) = |\{k \leq K_T | N_{t_k}^i(s^i, a^i) > 2N_{t_{k-1}}^i(s^i, a^i)\}|.$$

As a result, $\gamma^i(s^i, a^i) \leq \log N_{T+1}^i(s^i, a^i)$. Note that for any $i \in [n]$, $N_{T+1}^i(s^i, a^i) \leq T$ and we have $2S^i$ state-action pairs. Then, we have

$$\begin{aligned} \gamma &\leq 1 + \sum_{i \in [n]} \sum_{(s^i, a^i)} \gamma^i(s^i, a^i) \leq 1 + \sum_{i \in [n]} \sum_{(s^i, a^i)} \log N_{T+1}^i(s^i, a^i) \\ &= 1 + \sum_{i \in [n]} 2S^i \log T \leq 2\bar{S}_n \log T. \end{aligned}$$

B. Bound on $\mathcal{R}_1(T)$ (Lemma 2.2)

$\mathcal{R}_1(T)$ is a telescoping sum, which can be simplified as follows:

$$\begin{aligned} \mathcal{R}_1(T) &= \mathbb{E} \left[\sum_{k=1}^{K_T} \sum_{t=t_k}^{t_{k+1}-1} [\mathbf{V}_k(\mathbf{S}_{t+1}) - \mathbf{V}_k(\mathbf{S}_t)] \right] \\ &= \mathbb{E} \left[\sum_{k=1}^{K_T} [\mathbf{V}_k(\mathbf{S}_{t_{k+1}}) - \mathbf{V}_k(\mathbf{S}_{t_k})] \right] \leq 2 \frac{\alpha R_{\max}}{1 - \lambda^*} \mathbb{E}[K_T] \end{aligned}$$

where the last inequality uses Lemma 1. The result then follows by substituting the value of K_T from Lemma 9.

C. Bound on $\mathcal{R}_2(T)$ (Lemma 2.3)

1) Notation: For any arm $i \in [n]$, let $N_t^i(s^i, a^i, s_+^i)$ denote the number of times $(S_\tau, A_\tau, S_{\tau+1})$ is equal to (s^i, a^i, s_+^i) until time t . Let $\hat{P}_t^i(s_+^i | s^i, a^i) = N_t^i(s^i, a^i, s_+^i) / (1 \vee N_t^i(s^i, a^i))$ denote the empirical distribution based on observations up to time t . For a given $\delta \in (0, 1)$, we define

$$\epsilon_\delta^i(\ell) = \sqrt{\frac{2S^i \log(1/\delta)}{1 \vee \ell}}. \quad (13)$$

2) Some Preliminary Results: In this section, we state some preliminary properties.

Lemma 10: Let $p, q \in \Delta(\mathcal{S})$. Then, for any function $f : \mathcal{S} \rightarrow \mathbb{R}$, we have $|\langle f, p \rangle - \langle f, q \rangle| \leq 0.5 \text{span}(f) \|p - q\|_1$.

Proof: Let $\bar{f} = (\max f + \min f)/2$. Then

$$\begin{aligned} |\langle f, p \rangle - \langle f, q \rangle| &= |\langle f - \bar{f}, p - q \rangle| \leq \|f - \bar{f}\|_\infty |\langle \mathbf{1}, p - q \rangle| \\ &\leq \frac{1}{2} \text{span}(f) \|p - q\|_1. \end{aligned}$$

Lemma 11: Consider any arm i , episode k , $\delta \in (0, 1)$, $\ell > 1$, and state-action pair (s^i, a^i) . Define events $\mathcal{E}_\ell^i = \{N_{t_k}^i(s^i, a^i) = \ell\}$, $\mathcal{F}^i = \{\|P^i(\cdot | s^i, a^i) - \hat{P}_{t_k}^i(\cdot | s^i, a^i)\|_1 \leq \epsilon_\delta(N_{t_k}^i(s^i, a^i))\}$, and $\mathcal{F}_k^i = \{\|P_k^i(\cdot | s^i, a^i) - \hat{P}_{t_k}^i(\cdot | s^i, a^i)\|_1 \leq \epsilon_\delta(N_{t_k}^i(s^i, a^i))\}$. Then, we have

$$\mathbb{P}\left(\|P^i - \hat{P}_{t_k}^i(\cdot | s^i, a^i)\|_1 > \epsilon_\delta^i(\ell) \mid \mathcal{E}_\ell^i\right) \leq \delta.$$

The above inequality implies that

$$\mathbb{E}\left[\|P^i - \hat{P}_{t_k}^i(\cdot | s^i, a^i)\|_1\right] \leq \mathbb{E}[\epsilon_\delta^i(N_{t_k}^i(s^i, a^i)) | \mathcal{F}^i] + 2\delta.$$

A similar bound holds if (P^i, \mathcal{F}^i) is replaced by (P_k^i, \mathcal{F}_k^i) .

Proof: Given arm i , state s^i of the arm and action a^i chosen for the arm, we know from [61] that for any $\varepsilon > 0$, the L_1 -deviation of the true and the empirical distributions over \mathcal{S}^i with $N_{t_k}^i(s^i, a^i) = \ell$ samples is bounded by

$$\begin{aligned} \mathbb{P}\left(\|P^i(\cdot | s^i, a^i) - \hat{P}_{t_k}^i(\cdot | s^i, a^i)\|_1 \geq \varepsilon \mid \mathcal{E}_\ell^i\right) \\ \leq 2S^i \exp\left(-\frac{\ell\varepsilon^2}{2}\right) < \exp\left(S^i - \frac{\ell\varepsilon^2}{2}\right). \end{aligned}$$

Let $\delta = \exp(S^i - \ell\varepsilon^2/2)$. Note that, $S^i \geq 2$, therefore $S^i + \log(1/\delta) \leq S^i \log(1/\delta)$. Hence,

$$\mathbb{P}\left(\|P^i(\cdot | s^i, a^i) - \hat{P}_{t_k}^i(\cdot | s^i, a^i)\|_1 > \sqrt{\frac{2S^i \log(1/\delta)}{1 \vee \ell}} \mid \mathcal{E}_\ell^i\right) < \delta.$$

The next result is driven by showing that $P((\mathcal{F}^i)^c) \leq \delta$ and

$$\begin{aligned} \mathbb{E}\left[\|P^i(\cdot | s^i, a^i) - \hat{P}_{t_k}^i(\cdot | s^i, a^i)\|_1\right] \\ \leq 2\delta + \mathbb{E}[\epsilon_\delta^i(N_{t_k}^i(s^i, a^i)) | \mathcal{F}^i]. \end{aligned}$$

See [58] for details. \blacksquare

Lemma 12: For any episode k , and $\delta \in (0, 1)$, we have

$$\begin{aligned} \mathbb{E}\left[\|P_\star(\cdot | \mathbf{s}, \mathbf{a}) - P_k(\cdot | \mathbf{s}, \mathbf{a})\|_1\right] \\ \leq 4n\delta + \sum_{i \in [n]} \sum_{f \in \{\mathcal{F}^i, \mathcal{F}_k^i\}} \mathbb{E}[\epsilon_\delta^i(N_{t_k}^i(s^i, a^i)) | f]. \end{aligned}$$

Proof: The proof is as follows:

$$\begin{aligned} \mathbb{E}\left[\|P_\star(\cdot | \mathbf{s}, \mathbf{a}) - P_k(\cdot | \mathbf{s}, \mathbf{a})\|_1\right] \\ \stackrel{(a)}{=} \sum_{i \in [n]} \mathbb{E}\left[\|P^i(\cdot | s^i, a^i) - P_k^i(\cdot | s^i, a^i)\|_1\right] \\ \stackrel{(b)}{\leq} \sum_{i \in [n]} \mathbb{E}\left[\|P^i(\cdot | s^i, a^i) - \hat{P}_{t_k}^i(\cdot | s^i, a^i)\|_1\right] \\ + \|P_k^i(\cdot | s^i, a^i) - \hat{P}_{t_k}^i(\cdot | s^i, a^i)\|_1 \end{aligned}$$

where (a) follows from [35, Lemma 13] and (b) follows from triangle inequality. The result then follows from Lemma 11. \blacksquare

3) Bounding $\mathcal{R}_2(T)$: Now, consider the inner summation in the expression for $\mathcal{R}_2(T)$

$$\begin{aligned} \mathbb{E}\left[\langle P_k(\cdot | \mathbf{S}_t, \mathbf{A}_t), \mathbf{V}_k \rangle - \mathbf{V}_k(\mathbf{S}_{t+1})\right] \\ \stackrel{(a)}{\leq} \mathbb{E}\left[\frac{1}{2} \text{span}(\mathbf{V}_k) \|P_k(\cdot | \mathbf{S}_t, \mathbf{A}_t) - P_\star(\cdot | \mathbf{S}_t, \mathbf{A}_t)\|_1\right] \\ \stackrel{(b)}{\leq} \frac{\alpha R_{\max}}{1 - \lambda^*} \mathbb{E}\left[\|P_k(\cdot | \mathbf{S}_t, \mathbf{A}_t) - P_\star(\cdot | \mathbf{S}_t, \mathbf{A}_t)\|_1\right] \end{aligned} \quad (14)$$

where (a) follows from Lemma 10 and (b) follows from Lemma 1. Then, by Lemma 12, we have

$$\begin{aligned} \mathcal{R}_2(T) &= \mathbb{E}\left[\sum_{k=1}^{K_T} \sum_{t=t_k}^{t_{k+1}-1} [P_k \mathbf{V}_k](\mathbf{S}_t) - \mathbf{V}_k(\mathbf{S}_{t+1})\right] \\ &\leq \frac{\alpha R_{\max}}{1 - \lambda^*} \sum_{k=1}^{K_T} \sum_{t=t_k}^{t_{k+1}-1} \left(4n\delta \right. \\ &\quad \left. + \sum_{(i,f) \in \mathcal{D}_k} \mathbb{E}[\epsilon_\delta^i(N_{t_k}^i(s^i, a^i)) | f]\right) \end{aligned} \quad (15)$$

where $\mathcal{D}_k = \{(i, f) : i \in [n], f \in \{\mathcal{F}^i, \mathcal{F}_k^i\}\}$.

For the first inner term of (15), we have

$$\sum_{k=1}^{K_T} \sum_{t=t_k}^{t_{k+1}-1} 4n\delta = \sum_{t=1}^T 4n\delta = 4n\delta T. \quad (16)$$

For the second inner term of (15), fix $(i, f) \in \mathcal{D}_k$ and let $\bar{\delta} = 2\sqrt{S^i \log(1/\delta)}$. Note $\sum_{k=1}^{K_T} \sum_{t=t_k}^{t_{k+1}-1} = \sum_{t=1}^T$. Then,

$$\begin{aligned} \sum_{k=1}^{K_T} \sum_{t=t_k}^{t_{k+1}-1} \mathbb{E}\left[\frac{\bar{\delta}/\sqrt{2}}{\sqrt{1 \vee N_{t_k}^i(S_t^i, A_t^i)}} | f\right] \\ = \bar{\delta} \sum_{(s^i, a^i)} \sum_{t=1}^T \mathbb{E}\left[\mathbf{1}(S_t^i = s^i, A_t^i = a^i) \sqrt{\frac{1}{1 \vee N_t^i(s^i, a^i)}} | f\right] \end{aligned}$$

$$\begin{aligned}
&= \bar{\delta} \sum_{(s^i, a^i)} \mathbb{E} \left[\mathbb{1}(N_{T+1}^i(s^i, a^i) > 0) + \sum_{j=1}^{N_{T+1}^i(s^i, a^i)-1} \frac{1}{\sqrt{j}} \middle| f \right] \\
&\leq \bar{\delta} \sum_{(s^i, a^i)} \mathbb{E} \left[\mathbb{1}(N_{T+1}^i(s^i, a^i) > 0) + 2\sqrt{N_{T+1}^i(s^i, a^i)} \middle| f \right] \\
&\leq \bar{\delta} \sum_{(s^i, a^i)} 3\mathbb{E} \left[\sqrt{N_{T+1}^i(s^i, a^i)} \middle| f \right] \\
&\stackrel{(a)}{\leq} 3\bar{\delta} \mathbb{E} \left[\sqrt{2S^i \sum_{(s^i, a^i)} N_{T+1}^i(s^i, a^i)} \middle| f \right] \\
&\stackrel{(b)}{=} 6\sqrt{2}S^i \sqrt{T \log(1/\delta)} \tag{17}
\end{aligned}$$

where (a) uses the Cauchy–Schwartz inequality and (b) uses the fact that $\sum_{(x^i, a^i)} N_{T+1}^i(x^i, a^i) = T$. Adding (17) over $(i, f) \in \mathcal{D}_k$, we have

$$\begin{aligned}
&\sum_{k=1}^{K_T} \sum_{t=t_k}^{t_{k+1}-1} \sum_{(i, f) \in \mathcal{D}_k} \mathbb{E} [\epsilon_\delta^i(N_{t_k}^i(S_t^i, A_t^i)) | f] \\
&\leq 12\sqrt{2}\bar{S}_n \sqrt{T \log(1/\delta)}. \tag{18}
\end{aligned}$$

Finally, by setting $\delta = 1/T$, and Substituting (18), (17), and (16) in (15), we get the final result.

D. Bound on $\mathcal{R}_2(T)$ (Lemma 3)

1) Some Preliminary Results: For any arm $i \in [n]$, let $N_t^i(s^i, a^i, s_{+}^i)$ and $\hat{P}_t^i(s_{+}^i | s^i, a^i)$ denote the same variables, as defined in Section C.1.

Lemma 13: For any Lipschitz function $f : \mathcal{X} \rightarrow \mathbb{R}$ with Lipschitz coefficient L_f , and any probability measures ζ_1 and ζ_2 on (\mathcal{X}, d_X) we have

$$\left| \sum_{x \in \mathcal{X}} f(x) \zeta_1(x) - \sum_{x \in \mathcal{X}} f(x) \zeta_2(x) \right| \leq L_f \mathcal{K}(\zeta_1, \zeta_2).$$

Proof: The result is immediately derived from the definition of the Kantorovich distance. ■

Lemma 14 (From Theorem 2 of [62]): Let ν denote a probability measure on $(\mathbb{R}, |\cdot|)$ and let $\hat{\nu}_n$ denote the estimated probability measure by n samples from ν . Then, for all $n \geq 1$ and all $\epsilon > 0$, there exist constants C and c which depend on ν such that $\mathbb{P}(\mathcal{K}(\nu, \hat{\nu}_n) \geq \epsilon) \leq C \exp(-cn\epsilon) \mathbb{1}(\epsilon \leq 1) + C \exp(-cn\epsilon^2) \mathbb{1}(\epsilon > 1)$.

Proof: The lemma follows directly by applying [62, Th. 2] and setting $d = 1$, $p = 1$, and $\alpha = 2$ which satisfies [62, condition (C1)]. ■

Let $\epsilon_\delta(\ell) = \sqrt{\log(C/\delta)/(c(1 \vee \ell))}$.

Lemma 15: Consider any arm i , any episode k , $\delta \in (0, 1)$, $\ell > 1$, and state-action pair (s^i, a^i) . Define events $\mathcal{F}^i = \{\mathcal{K}(P^i, \hat{P}_{t_k}^i(\cdot | s^i, a^i)) \leq \epsilon_\delta(N_{t_k}^i(s^i, a^i))\}$, and $\mathcal{F}_k^i = \{\mathcal{K}(\tilde{P}_k^i, \hat{P}_{t_k}^i(\cdot | s^i, a^i)) \leq \epsilon_\delta(N_{t_k}^i(s^i, a^i))\}$. We have

$$\mathbb{P}(\mathcal{K}(P^i, \hat{P}_{t_k}^i(\cdot | s^i, a^i)) > \epsilon_\delta^i(\ell) \mid \mathcal{E}_\ell^i) \leq \delta.$$

Furthermore, the above inequality implies

$$\begin{aligned}
\mathbb{E}[\mathcal{K}(P^i, \hat{P}_{t_k}^i(\cdot | s^i, a^i))] &\leq \mathbb{E}[\epsilon_\delta^i(N_{t_k}^i(s^i, a^i)) | \mathcal{F}^i] \\
&\quad + 2 \text{diam}(S^i) \delta.
\end{aligned}$$

A similar bound holds if (P^i, \mathcal{F}^i) is replaced by (P_k^i, \mathcal{F}_k^i) .

Proof: The result follows by using Lemmas 13 and 14 and a similar approach as the proof of Lemma 11. ■

Lemma 16: For any episode k , and $\delta \in (0, 1)$, we have

$$\begin{aligned}
&\mathbb{E}[\mathcal{K}(P_\star(\cdot | s, a), P_k(\cdot | s, a))] \\
&\leq 4nD_{\max} \delta + \sum_{f \in \{\mathcal{F}^i, \mathcal{F}_k^i\}} \sum_{i \in [n]} \mathbb{E}[\epsilon_\delta^i(N_{t_k}^i(s^i, a^i)) \mid f].
\end{aligned}$$

Proof: The proof is similar to that of Lemma 12, where we use the Kantorovich distance instead of total variation distance. The equivalent of equality (a) (in the proof of Lemma 12) follows from [59, Lemma 4], and the rest of the argument follows from Lemma 15. ■

2) Bounding $\mathcal{R}_2(T)$: First, consider the inner summation in the expression for $\mathcal{R}_2(T)$

$$\begin{aligned}
&\mathbb{E}[\langle P_k(\cdot | S_t, \mathbf{A}_t), \mathbf{V}_k \rangle - \mathbf{V}_k(S_{t+1})] \\
&\stackrel{(a)}{\leq} L_v \mathbb{E}[\mathcal{K}(P_\star(\cdot | S_t, \mathbf{A}_t) - P_k(\cdot | S_t, \mathbf{A}_t))] \tag{19}
\end{aligned}$$

where (a) follows from Lemma 13. Then, by Lemma 16, we have

$$\begin{aligned}
\mathcal{R}_2(T) &= \mathbb{E} \left[\sum_{k=1}^{K_T} \sum_{t=t_k}^{t_{k+1}-1} [P_k \mathbf{V}_k](S_t) - \mathbf{V}_k(S_{t+1}) \right] \\
&\leq L_v \sum_{k=1}^{K_T} \sum_{t=t_k}^{t_{k+1}-1} \left(4nD_{\max} \delta \right. \\
&\quad \left. + \sum_{(i, f) \in \mathcal{D}_k} \mathbb{E}[\epsilon_\delta^i(N_{t_k}^i(S_t^i, A_t^i)) \mid f] \right) \tag{20}
\end{aligned}$$

where $\mathcal{D}_k = \{(i, f) : i \in [n], f \in \{\mathcal{F}^i, \mathcal{F}_k^i\}\}$.

For the first inner term of (20), we have

$$\sum_{k=1}^{K_T} \sum_{t=t_k}^{t_{k+1}-1} 4nD_{\max} \delta = \sum_{t=1}^T 4nD_{\max} \delta = 2nD_{\max} \delta T. \tag{21}$$

For the second inner term of (20), we can follow an argument similar to (18) to show that:

$$\begin{aligned}
&\sum_{k=1}^{K_T} \sum_{t=t_k}^{t_{k+1}-1} \sum_{(i, f) \in \mathcal{D}_k} \mathbb{E}[\epsilon_\delta^i(N_{t_k}^i(S_t^i, A_t^i)) \mid f] \\
&\leq 12\sqrt{\frac{\bar{S}_n \log(C/\delta) T}{c}}. \tag{22}
\end{aligned}$$

See [58] for details. Finally, by setting $\delta = 1/T$, and substituting (22) and (21) in (20), we get the result.

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